DISTRIBUTIVE MEREOTOPOLOGY: Extended Distributive Contact Lattices Tatyana Ivanova and Dimiter Vakarelov

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Abstract

The notion of contact algebra is one of the main tools in the region based theory of space. It is an extension of Boolean algebra with an additional relation C called contact. The elements of the Boolean algebra are considered as formal representations of spatial regions as analogs of physical bodies and Boolean operations are considered as operations for constructing new regions from given ones and also to define some mereological relations between regions as part-of, overlap and underlap. The contact relation is one of the basic mereotopological relations between regions expressing some topological nature. It is used also to define some other important mereotopological relations like non-tangential inclusion, dual contact, external contact and others. Most of these definitions are given by means of the operation of Boolean complementation. There are, however, some problems related to the motivation of the operation of Boolean complementation. In order to avoid these problems we propose a generalization of the notion of contact algebra by dropping the operation of complement and replacing the Boolean part of the definition by distributive lattice. First steps in this direction were made in [8, 9] presenting the notion of distributive contact lattice based on contact relation as the only mereotopological relation. In this paper we consider as nondefinable primitives the relations of contact, nontangential inclusion and dual contact, extending considerably the language of distributive contact lattices. Part I of the paper is devoted to a suitable axiomatization of the new language called extended distributive contact lattice (EDC-lattice) by means of universal first-order axioms true in all contact algebras. EDClattices may be considered also as an algebraic tool for certain subarea of mereotopology, called in this paper distributive mereotopology. The main result of Part I of the paper is a representation theorem, stating that each EDC-lattice can be isomorphically embedded into a contact algebra, showing in this way that the presented axiomatization preserves the meaning of mereotopological relations without considering Boolean complementation. Part II of the paper is devoted to topological representation theory of EDC-lattices, transferring into the distributive case important results from the topological representation theory of contact algebras. It is shown that under minor additional assumptions on distributive lattices as extensionality of the definable relations of overlap or underlap one can preserve the good topological interpretations of regions as regular closed or regular open sets in topological space.

Keywords: mereotopology, distributive mereotopology, contact algebras, distributive contact algebras, extended distributive contact algebras, topological representations.

1 Introduction

In this paper we continue the research line started in the publications [8, 9], devoted to certain non-classical approach to the region-based theory of space (RBTS), which roots goes back mainly to Whitehead [30]. In contrast to the classical Euclidean approach, in which the notion of point is taken as one of the basic primitive notions in geometry and geometric figures are considered as sets of points, RBTS adopts as primitives the more realistic spatial notion of region (as an abstraction of spatial or physical body), together with some basic relations and operations on regions. Some of these relations come from mereology (see [20]): e.g., part-of $(x \leq y)$, overlap (xOy), its dual underlap (xOy), and some others definable in terms of these. RBTS extends classical mereology by considering some new relations among regions which are topological in nature, such as contact (xCy), nontangential part-of $(x \ll y)$, dual contact (xCy), and some others definable by means of the contact and part-of relations. This is one of the reasons that the extension of mereology with these new relations is commonly called *mereotopology*. There is no clear difference in the literature between RBTS and mereotopology, and by some authors RBTS is related rather to the so called *mereogeometry*, while mereotopology is considered only as a kind of point-free topology, considering mainly topological properties of things. In this paper we consider all these names almost as synonyms representing collections of various point-free theories of space. According to Whitehead the point-free approach to space should not disregard points at all - on the contrary, they are suitable high level abstractions which, as such, should not be put on the base of the theory, but have to be definable by means of the other primitive notions of the theory. The Whitehead's criticism is based on the fact that points, as well as the other primitive notions in Euclidean geometry like lines and planes, do not have separate existence in reality, while for instance, spatial bodies as cubes, prisms, pyramids, balls, etc are things having analogs in reality. In this sense the point-free approach to space can be considered as certain equivalent re-formulation of the classical point-based approach by means of more realistic primitive notions.

Survey papers about RBTS (and mereotopology) are [24, 4, 14] (see also the handbook [1] and [3] for some logics of space). Let us mention that in a sense RBTS had been reinvented in computer science, because of its more simple way of representing qualitative spatial information and in fact it initiated a special field in Knowledge Representation (KR) called Qualitative Spatial Representation and Reasoning (QSRR). One of the most popular systems in QSRR is the Region Connection Calculus (RCC) introduced in [18]. Note that RCC influenced various investigations in the field both of theoretical and applied nature. Survey papers about applications of RBTS and mereotopology in various applied areas are, for instance, [5] and the book [16].

Let us note that one of the main algebraic tools in mereotopology is the notion of *contact algebra*, which appears in the literature under different names and formulations as extensions of Boolean algebra with some mereotopological relations [27, 21, 25, 26, 4, 10, 6, 7]. The simplest system, called just *contact algebra* was introduced in [6] as an extension of Boolean algebra $\underline{B} = (B, 0, 1, ..., +, *)$ with a binary relation C called *contact* and satisfying several simple axioms:

(C1) If aCb, then $a \neq 0$ and $b \neq 0$,

- (C2) If aCb and $a \leq a'$ and $b \leq b'$, then a'Cb',
- (C3) If aC(b+c), then aCb or aCc,
- (C4) If aCb, then bCa,
- (C5) If $a.b \neq 0$, then aCb.

The elements of the Boolean algebra are called regions and the Boolean operations can be considered as some constructions of new regions by means of given ones. In this definition Boolean algebra stands for the mereological component, while the contact relation C stands for the mereotopological component of the system. For instance the mereological relations overlap O, underlap (dual overlap) \widehat{O} and part-of \leq have the following definitions: $aOb \leftrightarrow_{def} a.b \neq 0$, $a\widehat{O}b \leftrightarrow_{def} a+b \neq 1$ and \leq is just the lattice ordering. The unite element 1 is the region containing as its parts all regions, and the zero region 0 symbolize the non-existing region and can be used to define the ontological predicate of existence: a exists $\leftrightarrow_{def} a \neq 0$. According to these definitions the axiom (C1) says that if a and b are in a contact then they exist, and axiom (C5) says that overlapping regions are in a contact.

By means of the contact relation one can define other mereotopological relations: dual contact $a\widehat{C}b \leftrightarrow_{def} a^*Cb^*$, non-tangential part-of $a \ll b \leftrightarrow_{def} a\overline{C}b^*$, and some others.

Intuitively if we consider regions as certain sets of points, then contact aCb means that a and b share a common point, part-of $a \le b$ means that all points of a are points of b, overlap aOb means that a and b share an existing region (just $a.b \ne 0$ is a part both of a and of b), underlap $a\widehat{O}b$ means that there exists a non-universal region containing both a and b (just $a+b\ne 1$ contains both a and b).

Let us note that standard model of Boolean algebra is the algebra of subsets of a given universe, so in such a model regions are pure sets and the mereological relations between regions are just the Boolean relations between sets. In this model one can not distinguish boundary and internal points of a given region and hence it can not express all kinds of contact, for instance, the so called external contact in which the contacting regions share only a boundary point (external contact is definable by the formula $aCb \wedge aOb$). For this reason standard point models of contact algebras are of topological nature and consist of the Boolean algebras of regular closed sets in a given topological space and the contact between two such sets means that they have a common point. Another topological model of contact algebra is the Boolean algebra of regular open sets of a topological space, but in this model contact is not so intuitive and is definable by the formula: $aCb \leftrightarrow_{def} Cl(a) \cap Cl(b) \neq \emptyset$, where Cl(a) is the topological closure operation. Let us mention that the topological representation theory of contact algebras can be treated just as a realization of the Whitehead's idea of defining points and of recreation the point-based structure of the corresponding kind of space within a point-free system (see, for instance, the surveys [24, 4]).

One of the motivations to put Boolean algebra on the base of the notion of contact algebra is based on the remark given by Tarski (see for this [20], page 25) that one of the most popular mereological systems, namely the system of Lesniewski, can be identified with the complete Boolean algebra with zero deleted. If we are not interested in infinite unions and intersections then we

can accept just Boolean algebra (with zero considered as non-existing region, as mentioned above). In the papers [8, 9] a generalization of the notion of contact algebra is presented just by replacing the Boolean algebra by means of a (bounded) distributive lattice and obtaining in this way the notion of distributive contact lattice. Some motivations for this generalization are the following. First, that Boolean algebra is a bounded distributive lattice and that the axioms of the contact relation do not use the operation of Boolean complementation * and have the same formulation in the language of bounded distributive lattice. Second, that the same can be said for the basic mereological relations partof, overlap and underlap - they have definitions in the language of distributive lattice without the operation of Boolean complement. Third, that the representation theory for distributive lattices is quite similar to the corresponding theory of Boolean algebras and we wanted to see if this can help us in transferring the topological representation theory of contact algebras to the more general theory of distributive contact lattices, keeping the topological meaning of regions as regular closed sets. And finally, one philosophical motivation: the meaning of the Boolean complementation a^* is not well motivated: if the region a represents a physical body, then what kind of body represents a^* ? In the point-based models this is "the rest out of a" from the "whole space", the latter identified with the sum of all observed regions, the unit region 1. However, if we extend the area of our observation we will obtain another unit, and then a^* will be changed. But it is natural to assume that physical bodies should not depend on the area of observation in which they are included. As a result of this generalization, one can see that the paper [9] generalizes almost all from the topological representation theory of contact algebras developed for instance in [10, 6] and even more; on the distributive case one can see some deep features which can not be observed in the Boolean case. For instance in the Boolean case mereological relations have some hidden properties which in the distributive case are not always fulfilled and have to be postulated explicitly (this is the so called *extensionality property* for the underlap and overlap relations). However, the obtained generalization in [8, 9] has some open problems. The mereotopological relations of non-tangential part-of and dual contact in contact algebras have definitions by means of the operation of complementation. However these relations have a meaning in topological representation of contact algebras which does not depend on the operation of complementation on regular closed sets. Namely, if a and b are regular closed subsets of a topological space X, then $a \ll b$ iff $a \subseteq Int(b)$ and aCb iff $Int(a) \cup Int(b) \neq X$, where Intis the topological operation of interior of a set. Thus, it will be interesting to add these relations as primitives to the language of distributive contact lattices and to axiomatize them by means of a set of universal first-order axioms and then to extend the topological representation theory from [9]. This is one of the main open problems in [9] which positive solution is subject of the present paper. One of the motivations for this extension of the language of distributive contact lattice is that in this way we obtain a system with full duality: contact C is dual to the dual contact \hat{C} and non-tangential part-of \ll is dual to it converse \gg and this symmetry makes possible to obtain proofs by duality. The obtained new algebraic mereotopological system is named Extended Distributive Contact Lattice, EDC-lattice for short. We will consider in the paper the topological representation theory of some axiomatic extensions of EDC-lattices with new axioms yielding representations in better topological spaces, generalizing in

this way the existing representation theory for contact algebras. Since all these investigations form a special subfield of mereotopology based on distributive lattices, we introduce for this subfield a special name - distributive mereotopology, which is included in the title of the present paper. Having in mind this terminology, then the subarea of mereotopology based on Boolean algebras should be named Boolean mereotopology. Similar special names for other subfields of mereotopology depending on the corresponding mereological parts also can be suggested: for instance the mereotopology considered in [15, 28, 29] is based on some non-distributive lattices - hence non-distributive mereotopology, and the mereotopological structures considered, for instance, in [17, 13] are pure relational and without any algebraic lattice-structure in the set of regions - hence relational mereotopology.

The paper is divided in two parts. Part I is devoted to the axiomatization of the three mereotopological relations of contact C, dual contact \widehat{C} and nontangential part-of \ll taken as primitives on the base of distributive lattice by means of universal first-order axioms, which remain true in contact algebras. The main result of this part is the abstract notion of Extended Distributive Contact Lattice (EDC-lattice) and an embedding theorem of EDC-latices into contact algebras, showing in this way that the meaning of the contact, dual contact and non-tangential part-of relations is preserved in the language of EDC-lattices. The method is based on a certain generalization of the Stone representation theory of distributive lattices [22, 2]. As a consequence of the embedding theorem one can consider EDC-lattice also as the universal fragment of contact algebra based on the signature of distributive lattice and mereotopological relations of contact C, dual contact \widehat{C} and non-tangential inclusion \ll . Relations of EDC-lattices with other mereotopological systems are also considered: EDC-lattices are relational mereotopological systems in the sense of [17], and the well known RCC-8 system of mereotopological relations is definable in the language of EDC-lattices.

Part II of the paper is devoted to the topological representation theory of EDC-lattices and some of their axiomatic extensions yielding representations in T_1 and T_2 spaces. Special attention is given to dual dense and dense representations (defined in Section 5.1) in contact algebras of regular closed and regular open subsets of topological spaces. The method is an extension of the representation theory of distributive contact lattices [9] and adaptation of some constructions from the representation theory of contact algebras [6, 7]. In the concluding Section we discuss some open problems and future plans with applications in qualitative spatial representation and reasoning.

PART I: EXTENDED DISTRIBUTIVE CONTACT LATTICES: AXIOMATIZATION AND EMBEDDING IN CONTACT ALGEBRAS

2 Extended distributive contact lattices. Choosing the right axioms

2.1 Contact algebras, distributive contact lattices and extended distributive contact lattices

As it was mention in the Introduction, contact algebra is a Boolean algebra $\underline{B} = (B, \leq, 0, 1, \cdot, +, *, C)$ with an additional binary relation C called contact, and satisfying the following axioms:

- (C1) If aCb, then $a \neq 0$ and $b \neq 0$,
- (C2) If aCb and $a \le a'$ and $b \le b'$, then a'Cb',
- (C3) If aC(b+c), then aCb or aCc,
- (C4) If aCb, then bCa,
- (C5) If $a.b \neq 0$, then aCb.

Let us note that on the base of (C4) we have (C3') (a + b)Cc implies aCc or bCc.

Remark 2.1 Observe that the above axioms are universal first-order conditions on the language of Boolean algebra with the C-relation and not containing the Boolean complementation *. This fact says that the axioms of C will be true in any distributive sublattice of \underline{B} . \square

The Remark 2.1 was one of the formal motivations for the definition of distributive contact lattice introduced in [8, 9]: the definition is obtained just by replacing the underlying Boolean algebra by a bounded distributive lattice $(D, \leq, 0, 1, +, \cdot)$ and taking for the relation C the same axioms. This makes possible to consider the main standard models of contact algebras, namely the algebras of regular closed or regular open sets of a topological space, also as the main models for distributive contact lattices, just by ignoring the Boolean complementation * in this models. This was guaranteed by Theorem 7 from [9] stating that every distributive contact lattice can be isomorphically embedded into a contact algebra, which fact indicates also that the choice of the set of axioms for distributive contact lattice is sufficient for proving this theorem. Since our main goal in the present paper is to obtain a definition of distributive contact lattice extended with relations of dual contact \widehat{C} and nontangential part-of \ll , we will follow here the above strategy, namely to choose universal firs-order statements for the relations C, \hat{C}, \ll as additional axioms which are true in arbitrary contact algebras and which guarantee the embedding into a contact algebra. The obtained algebraic system will be called extended distributive contact lattice. The next definition is a result of several preliminary experiments for fulfilling the above program.

Definition 2.2 Extended distributive contact lattice. Let $\underline{D} = (D, \leq ,0,1,+,\cdot,C,\widehat{C},\ll)$ be a bounded distributive lattice with three additional relations C,\widehat{C},\ll , called respectively **contact**, **dual contact** and **nontangential part-of**. The obtained system, denoted shortly by $\underline{D} = (D,C,\widehat{C},\ll)$, is called **extended distributive contact lattice** (EDC-lattice, for short) if it satisfies the axioms listed below.

Notations: if R is one of the relations \leq , C, \widehat{C} , \ll , then its complement is denoted by \overline{R} . We denote by \geq the converse relation of \leq and similarly \gg denotes the converse relation of \ll .

Axioms for C alone: The axioms (C1)-(C5) mentioned above.

Axioms for \widehat{C} alone:

- $(\widehat{C}1)$ If $a\widehat{C}b$, then $a, b \neq 1$,
- $(\widehat{C}2)$ If $a\widehat{C}b$ and $a' \leq a$ and $b' \leq b$, then $a'\widehat{C}b'$,
- $(\widehat{C}3)$ If $a\widehat{C}(b\cdot c)$, then $a\widehat{C}b$ or $a\widehat{C}c$,
- $(\widehat{C}4)$ If $a\widehat{C}b$, then $b\widehat{C}a$,
- $(\widehat{C}5)$ If $a+b \neq 1$, then $a\widehat{C}b$.

Axioms for \ll alone:

- $(\ll 1) \ 0 \ll 0$,
- $(\ll 2) \ 1 \ll 1$,
- $(\ll 3)$ If $a \ll b$, then $a \leq b$,
- $(\ll 4)$ If $a' \le a \ll b \le b'$, then $a' \ll b'$,
- $(\ll 5)$ If $a \ll c$ and $b \ll c$, then $(a+b) \ll c$,
- $(\ll 6)$ If $c \ll a$ and $c \ll b$, then $c \ll (a \cdot b)$,
- $(\ll 7)$ If $a \ll b$ and $(b \cdot c) \ll d$ and $c \ll (a+d)$, then $c \ll d$.

Mixed axioms:

- (MC1) If aCb and $a \ll c$, then $aC(b \cdot c)$,
- (MC2) If $a\overline{C}(b \cdot c)$ and aCb and $(a \cdot d)\overline{C}b$, then $d\widehat{C}c$,
- $(M\widehat{C}1)$ If $a\widehat{C}b$ and $c \ll a$, then $a\widehat{C}(b+c)$,
- $(M\widehat{C}2)$ If $a\overline{\widehat{C}}(b+c)$ and $a\widehat{C}b$ and $(a+d)\overline{\widehat{C}}b$, then dCc,
- $(M \ll 1)$ If $a\overline{\widehat{C}}b$ and $(a \cdot c) \ll b$, then $c \ll b$,
- $(M \ll 2)$ If $a\overline{C}b$ and $b \ll (a+c)$, then $b \ll c$.

Observation 2.3 Duality principle. For the language of EDCL we can introduce the following principle of duality: dual pairs $(0,1), (\cdot,+), (\leq, \geq), (C,\widehat{C}), (\ll, \gg)$. By means of these pairs for each statement (definition) A of the language we can define in an obvious way its dual \widehat{A} . Then by a routine verification one can see that for each axiom Ax from the list of axioms of EDCL its dual \widehat{Ax} is also true. On the base of this observation the proofs of dual statements will be omitted. Note, for instance, that each axiom from the first group (axioms for C

alone) is dually equivalent to the corresponding axiom from the second group (axioms for \widehat{C} alone) and vice versa, the third and fourth groups of axioms (axioms for \ll alone and mixed axioms) are closed under duality, for instance the axiom $(M\widehat{C}1)$ is dually equivalent to the axiom (MC1), and $(M \ll 2)$ is dually equivalent to $(M \ll 1)$. \square

2.2 Relational models of EDC-lattices

In order to prove that the axioms of EDC-lattices are true in contact algebras we will introduce a relational models of EDCL which are slight modifications of the relational models of contact algebras introduced in [7] and called there discrette contact algebras. The model is defined as follows.

Let (W, R) be a relational system where W is a nonempty set and R is a reflexive and symmetric relation in W and let a, b be arbitrary subsets of W. Define a contact relation between a and b as follows

(Def C_R) aC_Rb iff $\exists x \in a$ and $\exists y \in b$ such that xRy.

Then any Boolean algebra of subsets of W with thus defined contact is a contact algebra, and moreover, every contact algebra is isomorphic to a contact algebra of such a kind [7].

We will modify this model for EDCL as follows: instead of Boolean algebras of sets we consider only families of subsets containing the empty set \varnothing and the set W and closed under the set-union and set-intersection which are bounded distributive lattices of sets. Hence we interpret lattice constants and operations as follows: $0 = \varnothing$, 1 = W, $a \cdot b = a \cap b$, $a + b = a \cup b$. For the contact relation we preserve the definition (Def C_R). This modification is just a model of distributive contact lattice studied in [9].

Having in mind the definitions $a\widehat{C}b \leftrightarrow_{def} a^*Cb^*$ and $a \ll b \leftrightarrow_{def} a\overline{C}b^*$) in Boolean algebras, we introduce the following definitions for \widehat{C} and \ll (for some convenience we present the definition of the negation of \ll):

(Def
$$\widehat{C}_R$$
) $a\widehat{C}_R b$ iff $\exists x \notin a$ and $\exists y \notin b$ such that xRy , and (Def $\not \leqslant_R$) $a \not \leqslant_R b$ iff $\exists x \in a$ and $\exists y \notin b$ such that xRy .

Lemma 2.4 Let (W,R) be a relational system with reflexive and symmetric relation R and let \underline{D} be any collection of subsets of W which is a bounded distributive set-lattice with relations C, \widehat{C} and \ll defined as above. Then $(\underline{D}, C_R, \widehat{C}_R, \ll_R)$ is an EDC-lattice.

Proof. Routine verification that all axioms of EDC-lattice are true. \square EDC-lattice $\underline{D} = (D, C_R, \widehat{C}_R, \ll_R)$ over a relational system (W, R) will be called discrete EDC-lattice. If D is a set of all subsets of W then \underline{D} is called a full discrete EDC-lattice.

Corollary 2.5 The axioms of the relations C, \widehat{C} and \ll are true in contact algebras.

Proof. The proof follows by Lemma refRelationEDCL and the fact that every contact algebra can be isomorphically embedded into a discrete contact algebra over some relational system (W, R) wit reflexive and symmetric relation R [7]. \square

3 Embedding EDC-lattices into contact algebras

The main aim of this section is the proof a theorem stating that every EDC-lattice can be embedded into a full discrete EDC-lattice, which, of course is a Boolean contact algebra. As a consequence this will show that the axiomatization program for EDCL is fulfilled successfully. Since all axioms of EDC-lattice are universal first-order conditions, the axiomatization can be considered also as a characterization of the universal fragment of complement-free contact algebras based on the three relations. We will use in the representation theory a Stone like technique developed in [22] for the representation theory of distributive lattices.

3.1 Preliminary facts about filters and ideals in distributive lattices

We remaind some basic facts about filters and ideals in distributive lattices, for details see [2, 22].

Let \underline{D} be a distributive lattice. A subset F of D is called a filter in D if it satisfies the following conditions: (f1) $1 \in F$, (f2) if $a \in F$ and $a \leq b$ then $b \in F$, (f3) if $a, b \in F$ then $a.b \in F$. F is a proper filter if $0 \notin F$, F is a prime filter if it is a proper filter and $a + b \in F$ implies $a \in F$ or $b \in F$.

Dually, a subset I of D is an ideal if (i1) $0 \in I$, (i2) if $a \in I$ and $b \leq a$ then $b \in I$, (i3) if $a, b \in I$ then $a + b \in I$. I is a proper ideal if $1 \notin I$, I is a prime ideal if it is a proper ideal and $a.b \in I$ implies $a \in I$ or $b \in I$.

We will use later on some of the following facts without explicit mentioning.

Facts 3.1 Let \underline{D} be a bounded distributive lattice and Let F, F_1, F_2 be filters and I, I_1, I_2 be ideals.

- 1. The complement of a prime filter is a prime ideal and vice-versa.
- 2. $[a) = \{x \in D : a \le x\}$ is the smallest filter containing a;
 - $(a] = \{x \in D : x \leq a\}$ is the smallest ideal containing a.
- 3. $F_1 \oplus F_2 = \{c \in D : (\exists a \in F_1, b \in F_2)(a \cdot b \leq c)\} = \{a \cdot b : a \in F_1, b \in F_2\}$ is the smallest filter containing F_1 and F_2 .

$$[a) \oplus F = \{x \cdot y : a \le x, \ y \in F\}$$

 $I_1 \oplus I_2 = \{c \in D : (\exists a \in I_1, b \in I_2)(c \leq a+b)\} = \{a+b : a \in I_1, b \in I_2\}$ is the smallest ideal containing I_1 and I_2 .

$$(a] \oplus I = \{x + y : x \le a, y \in I\}.$$

In both cases the operation \oplus is associative and commutative.

4.
$$[a) \cap I = \emptyset$$
 iff $a \notin I$
If $(F \oplus [a)) \cap I \neq \emptyset$ then $(\exists x \in F)(a \cdot x \in I)$,
 $(a] \cap F = \emptyset$ iff $a \notin F$
If $F \cap (I \oplus (a]) \neq \emptyset$ then $(\exists x \in I)(a + x \in F)$.

The following three statements are well known in the representation theory of distributive lattices.

Lemma 3.2 Let F_0 be a filter, I_0 be an Ideal and $F_0 \cap I_0 = \emptyset$. Then:

- 1. Filter-extension Lemma. There exists a prime filter F such that $F_0 \subseteq F$ and $F \cap I_0 = \emptyset$.
- 2. Ideal-extension Lemma. There exists a prime ideal I such that $I_0 \subseteq I$ and $F_0 \cap I = \emptyset$.
- 3. Separation Lemma for filters and ideals. There exist a (prime) filter F and an (prime) ideal I such that $F_0 \subseteq F$, $I_0 \subseteq I$, $F \cap I = \emptyset$, and $F \cup I = D$.

Remark 3.3 Note that Filter-extension Lemma is dual to the Ideal-extension Lemma and that each of the three statement easily implies the other two. Normally they can be proved by application of the Zorn Lemma. The proof, for instance, of Filter-extension Lemma goes as follows. Apply the Zorn Lemma to the set $M = \{G : G \text{ is a filter}, F_0 \subseteq G \text{ and } G \cap I_0 = \emptyset\}$ and denote by F one of its maximal elements. Then it can be proved that F is a prime filter, and this finishes the proof. The sketched proof gives, however, an additional property of the filter F, namely

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(\forall x \notin F)(\exists y \in F)(x \cdot y \in I_0),
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which added to the formulation of the lemma makes it stronger. Since we will need later on this stronger version let us prove this property.

Suppose that $x \notin F$ and consider the filter $F \oplus [x)$. Since F is a maximal element of M, then $F \oplus [x)$ does not belong to M and consequently $F \oplus [x) \cap I_0 \neq \emptyset$. By the Fact 3.1, 4, there exists $y \in F$ such that $x \cdot y \in I_0$. We formulate this new statement below as $Strong\ ildeal$ -extension Lemma and its dual as $Strong\ ideal$ -extension Lemma. We do not know if these two statements for distributive lattices are new, but we will use them in the representation theorem in the next section. \square

Lemma 3.4 Let F_0 be a filter, I_0 be an Ideal and $F_0 \cap I_0 = \emptyset$. Then:

- 1. Strong filter-extension Lemma. There exists a prime filter F such that $F_0 \subseteq F$, $(\forall x \in F)(x \notin I_0)$ and $(\forall x \notin F)(\exists y \in F)(x \cdot y \in I_0)$.
- 2. Strong ideal-extension Lemma. There exists a prime ideal I such that $I_0 \subseteq I$, $(\forall x \in I)(x \notin F_0)$ and $(\forall x \notin I)(\exists y \in I)(x + y \in F_0)$.

3.2 Filters and Ideals in EDC-lattices

In the next two lemmas we list some constructions of filters and ideals in EDCL which will be used in the representation theory of EDC-lattices.

Lemma 3.5 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Then:

- 1. The set $I(x\overline{C}b) = \{x \in D : x\overline{C}b\}$ is an ideal,
- 2. the set $F(x\overline{\widehat{C}}b) = \{x \in D : x\overline{\widehat{C}}b\}$ is a filter,
- 3. the set $I(x \ll b) = \{x \in D : x \ll b\}$ is an ideal,
- 4. the set $F(x \gg b) = \{x \in D : x \gg b\}$ is a filter.

Proof. 1. By axiom (C1) $0\overline{C}b$, so $0 \in I(x\overline{C}b)$. Suppose $x \in I(x\overline{C}b)$ (hence $x\overline{C}b$) and $y \leq x$. Then by axiom (C2) $y\overline{C}b$). Let $x, y \in I(x\overline{C}b)$, hence $x\overline{C}b$ and $y\overline{C}b$. Then by axiom (C3) and (C4) we get $(x+y)\overline{C}b$ which shows that $x+y \in I(x\overline{C}b)$, which ends the proof of this case.

In a similar way one can proof 3. The cases 2. and 4. follow from 1. and 3. respectively by duality. \Box

Lemma 3.6 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice and Let Γ be a prime filter in D. Then:

- 1. The set $I(x\overline{C}\Gamma) = \{x \in D : (\exists y \in \Gamma)(x\overline{C}y)\}\$ is an ideal,
- 2. the set $F(x\overline{\widehat{C}}\overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x\overline{\widehat{C}}y)\}$ is a filter,
- 3. the set $I(x \ll \overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x \ll y)\}$ is an ideal,
- 4. the set $F(x \gg \Gamma) = \{x \in D : (\exists y \in \Gamma)(x \gg y)\}$ is a filter.

Proof. Note that the Lemma remains true if we replace Γ by a filter and $\overline{\Gamma}$ by an ideal.

1. The proof that $I(x\overline{C}\Gamma)$ satisfies the conditions (i1) and (i2) from the definition of ideal is easy. For the condition (i3) suppose $x_1, x_2 \in I(x\overline{C}\Gamma)$. Then $\exists y_1, y_2 \in \Gamma$ such that $x_1\overline{C}y_1$ and $x_2\overline{C}y_2$, Since Γ is a filter then $y = y_1 \cdot y_2 \in \Gamma$. Since $y \leq y_1$ and $y \leq y_2$, then by axiom (C2) we get $x_1\overline{C}y$ and $x_2\overline{C}y$. Then applying (C3') we obtain $(x_1 + x_2)\overline{C}y$, which shows that $x_1 + x_2 \in I(x\overline{C}\Gamma)$.

In a similar way one can prove 3. The proofs of 2 and 4 follow by duality from 1 and 3, taking into account that $\overline{\Gamma}$ is an ideal. \square

3.3 Relational representation theorem for EDC-lattices

Throughout this section we assume that $\underline{D}=(D,C,\widehat{C},\ll)$ is an EDC-lattice and let PF(D) and PI(D) denote the set of prime filters of \underline{D} and the set of prime ideals of D. Let $h(a)=\{\Gamma\in PF(D): a\in \Gamma\}$ be the well known Stone embedding mapping. We shall construct a canonical relational structure (W^c,R^c) related to \underline{D} putting $W^c=PF(D)$ and defining R^c for $\Gamma,\Delta\in PF(D)$ as follows:

 $\Gamma R^c \Delta \leftrightarrow_{def} (\forall a, b \in D)(a \in \Gamma, b \in \Delta \to aCb) \& (a \notin \Gamma, b \notin \Delta \to a\widehat{C}b) \& (a \in \Gamma, b \notin \Delta \to a \not\leqslant b) \& (a \notin \Gamma, b \in \Delta \to b \not\leqslant a)$

For some technical reasons and in order to use duality we introduce also the dual canonical structure $(\widehat{W}^c, \widehat{R}^c)$ putting $\widehat{W}^c = PI(D)$ and for $\Gamma, \Delta \in PI(D)$, $\Gamma \widehat{R}^c \Delta \leftrightarrow_{def} \overline{\Gamma} R^c \overline{\Delta}$.

Our aim is to show that the Stone mapping h is an embedding from \underline{D} into the EDC-lattice over (W^c, R^c) (see Section 2.4). First we need several technical lemmas.

Lemma 3.7 The canonical relations R^c and \hat{R}^c are reflexive and symmetric.

Proof. (For R^c) Symmetry is obvious by the definition of R^c and axioms (C4) and $(\widehat{C}4)$. In order to prove that $\Gamma R^c\Gamma$ suppose $a \in \Gamma$ and $b \in \Gamma$. Then $a \cdot b \in \Gamma$ and since Γ is a prime filter, then $a.b \neq 0$. Then by axiom (C5) we obtain aCb, which proves the first conjunct of the definition of R^c . For the

second conjunct suppose that $a \notin \Gamma$ and $b \notin \Gamma$, then, since Γ is a prime filter, $a+b \notin \Gamma$ and hence $a+b \neq 1$. Then by axiom $(\widehat{C}5)$ we get $a\widehat{C}b$. For the third conjunct suppose $a \in \Gamma$ and $b \notin \Gamma$, which implies that $a \nleq b$. Then by axiom $(\ll 3)$ we obtain $a \nleq b$. The proof of the last conjunct is similar.

(For
$$\widehat{R}^c$$
) - by duality. \square

Lemma 3.8 (i) aCb iff $(\exists \Gamma, \Delta \in PF(D))(a \in \Gamma \text{ and } b \in \Delta \text{ and } \Gamma R^c \Delta)$. (ii) $a \not\ll b$ iff $(\exists \Gamma, \Delta \in PF(D))(a \in \Gamma \text{ and } b \notin \Delta \text{ and } \Gamma R^c \Delta)$.

- **Proof.** (i) Note that the proof is quite technical, so we will present it with full details. The reasons for this are twofold: first to help the reader to follow it more easily, and second, to skip the details in a similar proofs.
 - (\Leftarrow) If $a \in \Gamma$ and $b \in \Delta$ then by the definition of R^c we obtain aCb.
 - (\Rightarrow) Suppose aCb.

The proof will go on several steps.

Step 1: construction of Γ . Consider the ideal $I(x\overline{C}b) = \{x \in D : x\overline{C}b\}$ (Lemma 3.5). Since aCb, $a \notin \{x \in D : x\overline{C}b\}$. Then $[a) \cap \{x \in D : x\overline{C}b\} = \varnothing$ and [a) is a filter (see Facts 3.1). By the Strong filter-extension lemma (see Lemma 3.4) there exists a prime filter Γ such that $[a) \subseteq \Gamma$ and $(\forall x \in \Gamma)(x \notin \{x \in D : x\overline{C}b\})$ and $(\forall x \notin \Gamma)(\exists y \in \Gamma)(x \cdot y \in \{x \in D : x\overline{C}b\})$. From here we conclude that Γ satisfies the following two properties:

```
(\#0) \ a \in \Gamma,
```

- (#1) If $x \in \Gamma$, then xCb, and
- (#2) If $x \notin \Gamma$, then there exists $y \in \Gamma$ such that $(x \cdot y)\overline{C}b$.

Step 2: construction of Δ . This will be done in two sub-steps.

Step 2.1 Consider the filters and ideals definable by Γ as in Lemma 3.6 $I(x\overline{C}\Gamma) = \{x \in D : (\exists y \in \Gamma)(x\overline{C}y)\}, F(x\overline{\widehat{C}}\Gamma) = \{x \in D : (\exists y \in \overline{\Gamma})(x\overline{\widehat{C}}y)\}, I(x \ll \overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x \ll y)\}, \text{ and } F(x \gg \Gamma) = \{x \in D : (\exists y \in \Gamma)(x \gg y)\}.$ In order to apply the Separation Lemma we will prove the following condition:

(#3)
$$F(x \gg \Gamma) \oplus F(x\overline{\widehat{C}}\overline{\Gamma}) \oplus [b) \cap I(x\overline{C}\Gamma) \oplus I(x \ll \overline{\Gamma}) = \emptyset$$
. Suppose that (#3) is not true, then for some $t \in D$ we have

- (1) $t \in F(x \gg \Gamma) \oplus F(x\widehat{C}\overline{\Gamma}) \oplus [b)$ and
- (2) $t \in I(x\overline{C}\Gamma) \oplus I(x \ll \overline{\Gamma}).$

It follows from (2) that $\exists k_1, k_2$ such that

- (3) $k_1 \in I(x \ll \overline{\Gamma})$ and
- (4) $k_2 \in I(x\overline{C}\Gamma)$ and
- $(5) t = k_1 + k_2.$

It follows from (1) that $\exists k_4, k_5, k_6 \in D$ such that

- (6) $k_4 \in F(x \gg \Gamma)$ and
- (7) $k_5 \in F(x\widehat{C}\overline{\Gamma})$ and
- (8) $k_6 \in [b)$ and
- $(9) t = k_4 \cdot k_5 \cdot k_6.$

From (5) and (9) we get

 $(10) k_1 + k_2 = k_4 \cdot k_5 \cdot k_6.$

It follows from (3), (4), (6) and (7) that

- (11) $\exists x_1 \in \overline{\Gamma}$ such that $k_1 \ll x_1$,
- (12) $\exists x_2 \in \Gamma \text{ such that } k_2 \overline{C} x_2$,

- (13) $\exists x_3 \in \Gamma$ such that $x_3 \ll k_4$,
- (14) $\exists x_4 \in \overline{\Gamma} \text{ such that } k_5 \overline{\widehat{C}} x_4.$

Let $x = x_1 + x_4$. Since $\overline{\Gamma}$ is an ideal, we obtain by (11) and (14) that

- (15) $x \in \overline{\Gamma}$ and $x \notin \Gamma$. Then by (#2) we get
- (16) $\exists y \in \Gamma$ such that $(x \cdot y)\overline{C}b$.

Let $z = x_2 \cdot x_3 \cdot y$. Then by (12), (13) and (16) we obtain that

(17) $z \in \Gamma$

and by (#1) that

(18)zCb.

From $x_1 \leq x$ and (11) by axiom ($\ll 4$) we get

(19) $k_1 \ll x$.

From $x_4 \leq x$ and (14) by axiom ($\widehat{C}2$) we obtain

(20) $k_5 \hat{C} x$.

From $z \leq x_2$ and (12) by axiom (C2) we get

(21) $k_2\overline{C}z$.

From $z \leq x_3$ and (13) by axiom ($\ll 4$) we obtain

(22) $z \ll k_4$.

We shall show that the following holds

(23) $z\overline{C}(b \cdot k_1)$.

Suppose for the sake of contradiction that $zC(b \cdot k_1)$. From $b \cdot k_1 \leq k_1$ and (19) by axiom ($\ll 4$) we get $(b \cdot k_1) \ll x$. From this fact and $zC(b \cdot k_1)$ by axiom (MC1) we obtain $(b \cdot k_1)C(z \cdot x)$. But we also have $b \cdot k_1 \leq b$, $z \cdot x \leq y \cdot x$, so by axiom (C2) we get $bC(y \cdot x)$ - a contradiction with (16).

The following condition holds

(24) $z\overline{C}(b \cdot k_2)$.

To prove this suppose for the sake of contradiction that $zC(b \cdot k_2)$. We also have $b \cdot k_2 \leq k_2$, so by axiom (C2) we get zCk_2 - a contradiction with (21).

Suppose that $zC(b \cdot (k_1 + k_2))$. By axiom (C3) we have $zC(b \cdot k_1)$ or $zC(b \cdot k_2)$ - a contradiction with (23) and (24). Consequently $z\overline{C}(b \cdot (k_1 + k_2))$ and by (10) we obtain $z\overline{C}(b \cdot k_4 \cdot k_5 \cdot k_6)$. But $b \leq k_6$ (from (8)), so $b \cdot k_4 \cdot k_5 \cdot k_6 = b \cdot k_4 \cdot k_5$. Consequently

(25) $z\overline{C}(b \cdot k_4 \cdot k_5)$.

From (18) and (22) by axiom (MC1) we get

 $(26) zC(b \cdot k_4).$

We shall show that the following condition holds

 $(27) (z \cdot x)C(b \cdot k_4)$

For to prove this suppose the contrary $(z \cdot x)C(b \cdot k_4)$. We also have $z \cdot x \leq y \cdot x$, $b \cdot k_4 \leq b$, so by axiom (C2) we get $(y \cdot x)Cb$ - a contradiction with (16).

From (25), (26) and (27) by axiom (MC2) we obtain $x\widehat{C}k_5$ - a contradiction with (20). Consequently (#3) is true.

Step 2.2: the construction of Δ . Applying the Filter extension Lemma to (#3) we obtain a prime filter Δ (and this is just the required Δ) such that:

1.
$$F(x \gg \Gamma) = \{x \in D : (\exists y \in \Gamma)(x \gg y\} \subseteq \Delta,$$

2.
$$F(x\overline{\widehat{C}}\overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x\overline{\widehat{C}}y)\} \subseteq \Delta,$$

3. $b \in \Delta$,

4.
$$I(x\overline{C}\Gamma) = \{x \in D : (\exists y \in \Gamma)(x\overline{C}y)\} \cap \Delta = \emptyset$$
,

5. $I(x \ll \overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x \ll y)\} \cap \Delta = \emptyset.$

Step 3: proof of $\Gamma R^c \Delta$. We will verify the four cases of the definition of R^c .

- Case 1: $y \in \Gamma$ and $x \in \Delta$. We have to show yCx. Suppose $y\overline{C}x$. Then $x\overline{C}y$ and by $y \in \Gamma$ we get $x \in I(x\overline{C}\Gamma)$. Then by 4. $x \notin \Delta$ a contradiction, hence yCx.
- Case 2: $y \in \Gamma$ and $x \notin \Delta$. Suppose $y \ll x$. Then $x \gg y$ and $y \in \Gamma$ implies $x \in F(x \gg \Gamma)$. By (1) $x \in \Delta$ a contradiction, hence $y \not\ll x$.
- Case 3: $y \notin \Gamma$ and $x \in \Delta$. Suppose $x \ll y$. Then $x \in I(x \ll \overline{\Gamma})$ and by 5. $x \notin \Delta$ a contradiction. Hence $x \not \ll y$.
- Case 4: $y \notin \Gamma$ and $x \notin \Delta$. Suppose $y\overline{\widehat{C}}x$. Then $x\overline{\widehat{C}}y$ and by 2. we obtain $x \in \Delta$ a contradiction. Hence $y\overline{\widehat{C}}x$.

Thus we have constructed prime filters Γ and Δ such that: $a \in \Gamma$, $b \in \Delta$ (item 3 from Step 2.2) and $\Gamma R^c \Delta$ (Step 3).

Proof of (ii). (\Leftarrow) If $a \in \Gamma$ and $b \notin \Delta$ then by the definition of R^c we obtain $a \not\ll b$.

 (\Rightarrow) Suppose $a \not\ll b$. The proof, as in (i), will go on several steps.

Step 1: construction of Γ . Consider the ideal $I(x \ll b) = \{x \in D : x \ll b\}$ (Lemma 3.5).

Since $a \not \ll b$, $a \not \in \{x \in D : x \ll b\}$. Then $[a) \cap \{x \in D : x \ll b\} = \emptyset$ and [a) is a filter (see FACTS 3.1). By the Strong filter-extension lemma (Lemma 3.4) there exists a prime filter Γ such that $[a) \subseteq \Gamma$ and $(\forall x \in \Gamma)(x \not \in \{x \in D : x \ll b\})$ and $(\forall x \not \in \Gamma)(\exists y \in \Gamma)(x \cdot y \in \{x \in D : x \ll b\})$. From here we conclude that Γ satisfies the following properties:

- $(\#0) \ a \in \Gamma,$
- (#1) If $x \in \Gamma$, then $x \not\ll b$, and
- (#2) If $x \notin \Gamma$, then there exists $y \in \Gamma$ such that $(x \cdot y) \ll b$.

Step 2: construction of Δ . This will be done in two sub-steps.

Step 2.1 Consider the filters and ideals definable by Γ as in Lemma 3.6 $I(x\overline{C}\Gamma) = \{x \in D : (\exists y \in \Gamma)(x\overline{C}y)\}, F(x\overline{\widehat{C}}\Gamma) = \{x \in D : (\exists y \in \overline{\Gamma})(x\overline{\widehat{C}}y)\}, I(x \ll \overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x \ll y)\}, \text{ and } F(x \gg \Gamma) = \{x \in D : (\exists y \in \Gamma)(x \gg y)\}.$ In order to apply the Filter-extension Lemma (Lemma 3.2) we will prove the following condition:

(#3) $F(x \gg \Gamma) \oplus F(x\overline{\widehat{C}}\overline{\Gamma}) \cap I(x \ll \overline{\Gamma}) \oplus I(x\overline{C}\Gamma) \oplus (b] = \emptyset$ Suppose that (#3) is not true. Consequently $\exists t$ such that

- (1) $t = k_1 \cdot k_2 = k_4 + k_5 + k_6$ for some $k_1, k_2, k_4, k_5, k_6 \in D$ and
- (2) $\exists x_1 \in \Gamma \text{ such that } x_1 \ll k_1$,
- (3) $\exists x_2 \in \overline{\Gamma} \text{ such that } k_2 \widehat{C} x_2$,
- (4) $\exists x_3 \in \overline{\Gamma}$ such that $k_4 \ll x_3$,
- (5) $\exists x_4 \in \Gamma \text{ such that } k_5 \overline{C} x_4$,
- (6) $k_6 \leq b$.

Let $z = x_2 + x_3$. Then by (3) and (4) we obtain $z \in \overline{\Gamma}$. By axiom ($\widehat{C}2$) we get

(7) $k_2\overline{\widehat{C}}z$.

By (4) and axiom ($\ll 4$) we get

(8) $k_4 \ll z$.

By $z \notin \Gamma$ and (#2) we have

(9) $\exists y \in \Gamma$ such that $(z \cdot y) \ll b$.

Let $x = x_1 \cdot x_4 \cdot y \cdot a$. Then by (#0), (2), (5) and (9) we get $x \in \Gamma$. By axiom $(\ll 4)$ we get

(10) $x \ll k_1$.

By (5), $x \le x_4$ and axiom (C2) we get

(11) $k_5\overline{C}x$.

From $x \in \Gamma$ by (#1) we obtain

(12) $x \not \ll b$.

From (10) by axiom ($\ll 4$) we get

(13) $x \ll (b + k_1)$

From (7) by axiom ($\widehat{C}2$) we obtain

 $(14) z\widehat{C}(b+k_2).$

From (9) by axiom ($\ll 4$) we get

 $(15) (z \cdot y) \ll (b + k_2).$

From (14) and (15) by axiom $(M \ll 1)$ we obtain $y \ll (b+k_2)$. We also have $x \leq y$ and by axiom $(\ll 4)$ we get

(16) $x \ll (b + k_2)$.

From (13) and (16) by axiom ($\ll 6$) we get $x \ll (b+k_1) \cdot (b+k_2)$. We have $(b+k_1) \cdot (b+k_2) = b+k_1 \cdot k_2 = b+k_4+k_5+k_6 = b+k_4+k_5$ (since $k_6 \leq b$ from (6)). Thus:

(17) $x \ll (b + k_4 + k_5)$.

Suppose (in order to obtain a contradiction) that $x \ll (b+k_4)$. From (9) and $x \cdot z \leq z \cdot y$ (which follows from the definitions of x and z) by axiom ($\ll 4$) we obtain $(x \cdot z) \ll b$. Using this fact, (8), $x \ll (b+k_4)$ and axiom ($\ll 7$) we get $x \ll b$ - a contradiction with (12). Consequently (18) $x \not\ll (b+k_4)$.

From (11) and (17) by axiom $(M \ll 2)$ we obtain $x \ll (b + k_4)$ - a contradiction with (18). Consequently (#3) is true.

Step 2.2: the construction of Δ . Applying the Filter-extension Lemma to (#3) we obtain a prime filter Δ (and this is just the required Δ) such that:

- 1. $F(x \gg \Gamma) = \{x \in D : (\exists y \in \Gamma)(x \gg y\} \subseteq \Delta,$
- 2. $F(x\overline{\widehat{C}}\overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x\overline{\widehat{C}}y)\} \subseteq \Delta,$
- 3. $b \notin \Delta$,
- 4. $I(x\overline{C}\Gamma) = \{x \in D : (\exists y \in \Gamma)(x\overline{C}y)\} \cap \Delta = \emptyset$,
- 5. $I(x \ll \overline{\Gamma}) = \{x \in D : (\exists y \in \overline{\Gamma})(x \ll y)\} \cap \Delta = \emptyset.$

Step 3: proof of $\Gamma R^c \Delta$. The proof is the same as in the corresponding step in (i).

To conclude: we have constructed prime filters Γ, Δ such that $\Gamma R^c \Delta$, $a \in \Gamma$ and $b \notin \Delta$, which finishes the proof of the lemma. \square

- **Lemma 3.9** (i) $a\widehat{C}b$ iff $(\exists \Gamma, \Delta \in PI(D))(a \in \Gamma \text{ and } b \in \Delta \text{ and } \Gamma \widehat{R}^c \Delta)$.
 - (ii) $a\widehat{C}b$ iff $(\exists \Gamma, \Delta \in PF(D))(a \notin \Gamma \text{ and } b \notin \Delta \text{ and } \Gamma R^c \Delta)$.
 - (iii) $a \gg b$ iff $(\exists \Gamma, \Delta \in PI(D))(a \in \Gamma \text{ and } b \notin \Delta \text{ and } \Gamma \widehat{R}^c \Delta)$.
 - (iv) $a \gg b$ iff $(\exists \Gamma, \Delta \in PF(D))(a \notin \Gamma \text{ and } b \in \Delta \text{ and } \Gamma R^c \Delta)$.
- **Proof.** (i) by duality from Lemma 3.8. Note that in this case Strong ideal-extension Lemma is used. The proof can follow in a "dual way" the steps of the proof of Lemma 3.8 (i).
 - (ii) is a corollary from (i).
 - (iii) by duality from Lemma 3.8 (ii) with the same remark as above.
 - (iv) is a corollary from (iii).□
- **Lemma 3.10** Let (W^c, R^c) be the canonical structure of $\underline{D} = (D, C, \widehat{C}, \ll)$ and $h(a) = \{U \in PF(D) : a \in U\}$ be the Stone mapping from D into the distributive lattice of all subsets of W^c . Then h is an embedding of \underline{D} into the EDC-lattice over (W^c, R^c) .
- **Proof.** It is a well known fact that h is an embedding of distributive lattice into the distributive lattice of all subsets of the set of prime filters PF(D) (see, [22, 2]). The only thing which have to be done is to show the following equivalences for all $a, b \in D$:
 - (i) aCb iff $h(a)C_{R^c}h(b)$,
 - (ii) $a\widehat{C}b$ iff $h(a)\widehat{C}_{R^c}h(b)$
 - (iii) $a \ll b$ iff $h(a) \ll_{R^c} h(b)$.

Note that these equivalences are another equivalent reformulation of Lemma 3.8 (i) and (ii) and Lemma 3.9 (ii) and (iv). \Box

Theorem 3.11 Relational representation Theorem of EDC-latices. Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Then there is a relational system $\underline{W} = (W, R)$ with reflexive and symmetric R and an embedding h into the EDC-lattice of all subsets of W.

Proof. The theorem is a corollary of Lemma $3.10.\Box$

Corollary 3.12 Every EDC-lattice can be isomorphically embedded into a contact algebra.

Proof. Since the lattice of all subsets of a given set is a Boolean algebra, then this is a corollary of Theorem 3.11. \Box

The following theorem states that the axiom system of EDC-lattice can be considered as an axiomatization of the universal fragment of contact algebras in the language of EDC-lattices.

Theorem 3.13 Let A be an universal first-order formula in the language of EDC-lattices. Then A is a consequence from the axioms of EDC-lattice iff A is true in all contact algebras.

Proof. The proof is a consequence from Corollary 3.12 and the fact that all axioms of EDC-lattice are universal first-order conditions and that A is also an universal first-order condition.

4 Relations to other mereotopologies

In this section we will compare EDC-lattices with other two mereotopologies: the *relational mereotopology* and *RCC-8*.

4.1 Relational mereotopology

Relational mereotopology is based on mereotopological structures introduced in [17]. These are relational structures in the form $(W, \leq, O, \widehat{O}, \ll, C, \widehat{C})$ axiomatizing the basic mereological relations part-of \leq , overlap O and dual overlap (underlap) \widehat{O} , and the basic mereotopological relations non-tangential part-of \ll , contact C and dual contact \widehat{C} . These relations satisfy the following list of universal first-order axioms:

diversal first-order axioms:
$$(\leq 0) \quad a \leq b \text{ and } b \leq a \rightarrow a = b \quad (\leq 1) \quad a \leq a,$$

$$(\leq 2) \quad a \leq b \text{ and } b \leq c \rightarrow a \leq c$$

$$(O1) \quad aOb \rightarrow bOa \quad (\widehat{O1}) \quad a\widehat{O}b \rightarrow b\widehat{O}a \quad (\widehat{O2}) \quad a\widehat{O}b \rightarrow a\widehat{O}a \quad (\widehat{O2}) \quad a\widehat{O}b \rightarrow a\widehat{O}a \quad (\widehat{O2}) \quad a\widehat{O}b \rightarrow a\widehat{O}a \quad (\widehat{O2}) \quad a\widehat{O}b \rightarrow a \leq b \quad (\widehat{O2}) \quad a\widehat{O}b \rightarrow a\widehat{O}b \rightarrow a\widehat{O}b \quad (\widehat{O3}) \quad a\widehat{O}b \rightarrow a\widehat{O}b \rightarrow a\widehat{O}b \quad (\widehat{O3}) \quad a\widehat{O}b \rightarrow a\widehat{O}b \rightarrow a\widehat{O}b \quad (\widehat{O3}) \quad a\widehat{O}b \rightarrow a\widehat{O}b \rightarrow a\widehat{O}b \quad (\widehat{C3}) \quad a\widehat{O}b$$

Note that all axioms of mereotopological structures are universal first-order conditions which are true in contact algebras under the standard definitions of the three basic mereological relations.

It is proved in [17] that each mereotopological structure is embeddable into a contact algebra (Theorem 26).

The following theorem relates EDC-lattices to mereotopological structures.

Theorem 4.1 Every EDC-lattice is a mereotopological structure under the standard definitions of the basic mereological relations.

Proof. Since all axioms of mereotopological structures are universal first-order sentences true in all contact algebras, then the statement follows from Theorem 3.13.

4.2 RCC-8 spatial relations

One of the most popular systems of topological relations in the community of QSRR is RCC-8. The system RCC-8 was introduced for the first time in [11]. It consists of 8 relations between non-empty regular closed subsets of arbitrary topological space. Having in mind the topological representation of contact algebras, it was given in [24] an equivalent definition of RCC-8 in the language of contact algebras:

Definition 4.2 The system RCC-8.

- disconnected $\mathbf{DC}(a, b)$: $a\overline{C}b$,
- external contact $-\mathbf{EC}(a,b)$: aCb and $a\overline{O}b$,
- partial overlap PO(a, b): aOb and $a \leq b$ and $b \leq a$,
- tangential proper part TPP(a,b): $a \le b$ and $a \not\ll b$ and $b \not\leq a$,
- tangential proper part⁻¹ TPP⁻¹(a,b): $b \le a$ and $b \not\ll a$ and $a \not\leq b$,
- nontangential proper part NTPP(a,b): $a \ll b$ and $a \neq b$,
- nontangential proper part⁻¹ NTPP⁻¹(a, b): $b \ll a$ and $a \neq b$,
- equal $\mathbf{EQ}(a, b)$: a = b.

Looking at this definition it can be easily seen that the RCC-8 relations are expressible in the language of EDC-lattices. Let us note that RCC-8 relations are not expressible in the language of distributive contact algebras from [9].

5 Additional axioms

In this Section we will formulate several additional axioms for EDC-lattices which are adaptations for the language of EDC-lattices of some known axioms considered in the context of contact algebras. First we will formulate some new lattice axioms for EDC-lattices - the so called extensionality axioms for the definable predicates of overlap - $aOb \leftrightarrow_{def} a \cdot b \neq 0$ and underlap - $a\widehat{O}b \leftrightarrow_{def} a + b \neq 1$.

(Ext O)
$$a \leq b \rightarrow (\exists c)(a \cdot c \neq 0 \text{ and } b \cdot c = 0)$$
 - extensionality of overlap,

(Ext
$$\widehat{O}$$
) $a \leq b \rightarrow (\exists c)(a+c=1 \text{ and } b+c\neq 1)$ - extensionality of underlap.

We say that a lattice is O-extensional if it satisfies (Ext O) and U-extensional if it satisfies (Ext \widehat{O}). Note that the conditions (Ext O) and (Ext \widehat{O}) are true in Boolean algebras but not always are true in distributive lattices (see [9] for some examples, references and additional information about these axioms).

We will study also the following extensionality axioms.

(Ext C)
$$a \neq 1 \rightarrow (\exists b \neq 0)(a\overline{C}b)$$
 - C-extensionality,

(Ext
$$\widehat{C}$$
) $a \neq 0 \rightarrow (\exists b \neq 1)(a\overline{\widehat{C}}b)$ - \widehat{C} -extensionality.

In contact algebras these two axioms are equivalent. It is proved in [9] that (Ext \widehat{O}) implies that (Ext C) is equivalent to the following extensionality principle considered by Whitehead [30]

(EXT C)
$$a \leq b \rightarrow (\exists c)(aCc \text{ and } b\overline{C}c)$$
.

Just in a dual way one can show that (Ext O) implies that (Ext \widehat{C}) is equivalent to the following condition

(EXT
$$\widehat{C}$$
) $a \not\leq b \to (\exists c)(b\widehat{C}c \text{ and } a\overline{\widehat{C}}c)$.

Let us note that (EXT C) and (EXT \widehat{C}) are equivalent in contact algebras. (Con C) $a \neq 0, b \neq 0$ and $a + b = 1 \rightarrow aCb$ - C-connectedness axiom and (Con \widehat{C}) $a \neq 1, b \neq 1$ and $a \cdot b = 0 \rightarrow a\widehat{C}b$ - \widehat{C} -connectedness axiom .

In contact algebras these axioms are equivalent and guarantee topological representation in connected topological spaces.

(Nor 1)
$$a\overline{C}b \to (\exists c, d)(c + d = 1, a\overline{C}c \text{ and } b\overline{C}d)$$
,

(Nor 2)
$$a\overline{\widehat{C}}b \to (\exists c, d)(c \cdot b = 0, a\overline{\widehat{C}}c \text{ and } b\overline{\widehat{C}}d)$$
,

(Nor 3)
$$a \ll b \to (\exists c)(a \ll c \ll b)$$
.

Let us note that the above three axioms are equivalent in contact algebras and are known by different names. For instance (Nor 1) comes from the proximity theory [23] as $Efremovich\ axiom$, (Nor 3) sometimes is called $interpolation\ axiom$. We adopt the name $normality\ axioms$ for (Nor 1), (Nor 2) and (Nor 3) because in topological representations they imply some normality conditions in the corresponding topological spaces. It is proved in [7] that (Nor 1) is true in the relational models (W,R) (see Section 2.2) if and only if the relation R is transitive and that (Nor 1) implies representation theorem in transitive models. In the next lemma we shall prove similar result using all normality axioms.

Lemma 5.1 Transitivity lemma. Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be a EDC-lattice satisfying the axioms (Nor1), (Nor 2) and (Nor 3) and let (W^c, R^c) be the canonical structure of D (see Section 3.3) Then:

- (i) R^c is a transitive relation.
- (ii) \underline{D} is representable in EDC-lattice over some system (W,R) with an equivalence relation R.

Proof. (i) Let Γ, Δ and Θ be prime filters in D such that

- (1) $\Gamma R^c \Delta$ and
- (2) $\Delta R^c \Theta$

and suppose for the sake of contradiction that

(3) $\Gamma \overline{R}^c \Theta$. By the definition of R^c we have to consider four cases.

Case 1: $\exists a \in \Gamma, b \in \Theta$ such that $a\overline{C}b$.

Then by (Nor 1) there exists c,d such that c+d=1, $a\overline{C}c$ and $b\overline{C}d$. Since c+d=1 then either $c\in\Delta$ or $d\in\Delta$. The case $c\in\Delta$ together with $a\in\Gamma$ imply by (1) aCc - a contradiction. The case $d\in\Delta$ together with $b\in\Theta$ imply by (2) bCd - again a contradiction.

Case 2: $\exists a \in \Gamma, b \notin \Theta$ such that $a \ll b$.

Then by (Nor 3) $\exists c$ such that $a \ll c$ and $c \ll b$. Consider the case $c \notin \Delta$. Then $a \in \Gamma$ and (1) imply $a \not\ll c$ a contradiction. Consider now $c \in \Delta$. Then $b \notin \Theta$ imply $c \not\ll b$ - again a contradiction.

In a similar way one can obtain a contradiction in the remaining two cases: Case 3: $\exists a \notin \Gamma, b \in \Theta$ such that $b \ll a$ and

Case 4: $\exists a \notin \Gamma, b \notin \Theta$ such that $b\widehat{C}a$.

(ii) The proof follows from (i) analogous to the proof of Theorem 3.11.□

Another kind of axioms which will be used in the topological representation theory in PART II are the so called rich axioms.

(U-rich
$$\ll$$
) $a \ll b \rightarrow (\exists c)(b+c=1 \text{ and } a\overline{C}c)$,

(U-rich
$$\widehat{C}$$
) $a\overline{\widehat{C}}b \to (\exists c, d)(a+c=1, b+d=1 \text{ and } c\overline{C}d)$.

(O-rich
$$\ll$$
) $a \ll b \to (\exists c)(a \cdot c = 0 \text{ and } c\overline{\widehat{C}}b),$

(O-rich C)
$$a\overline{C}b \to (\exists c, d)(a \cdot c = 0, b \cdot d = 0 \text{ and } c\overline{\widehat{C}}d).$$

Let us note that U-rich axioms will be used always with the U-extensionality axiom and that O-rich axioms will be used always with O-extensionality axiom. The following lemma is obvious.

Lemma 5.2 The axioms (U-rich \ll), (U-rich \widehat{C}), (O-rich \ll) and (O-rich C) are true in all contact algebras.

5.1 Some good embedding properties

Let $(D_1, C_1, \widehat{C}_1, \ll_1)$ and $(D_2, C_2, \widehat{C}_2, \ll_2)$ be two EDC-lattices. We will write $D_1 \leq D_2$ if D_1 is a substructure of D_2 , i.e., D_1 is a sublattice of D_2 , and the relations $C_1, \widehat{C}_1, \ll_1$ are restrictions of the relations $C_2, \widehat{C}_2, \ll_2$ on D_1 . Since we want to prove embedding theorems, it is valuable to know under what conditions we have equivalences of the form:

 D_1 satisfies some additional axiom iff D_2 satisfies the same axiom.

Definition 5.3 Dense and dual dense sublattice. Let D_1 be a distributive sublattice of D_2 . D_1 is called a dense sublattice of D_2 if the following condition is satisfied:

(Dense)
$$(\forall a_2 \in D_2)(a_2 \neq 0 \Rightarrow (\exists a_1 \in D_1)(a_1 \leq a_2 \text{ and } a_1 \neq 0)).$$

If h is an embedding of the lattice D_1 into the lattice D_2 then we say that h is a dense embedding if the sublattice $h(D_1)$ is a dense sublattice of D_2 .

Dually, D_1 is called a dual dense sublattice of D_2 if the following condition is satisfied:

(Dual dense)
$$(\forall a_2 \in D_2)(a_2 \neq 1 \Rightarrow (\exists a_1 \in D_1)(a_2 \leq a_1 \text{ and } a_1 \neq 1)).$$

If h is an embedding of the lattice D_1 into the lattice D_2 then we say that h is a Dual dense embedding if the sublattice $h(D_1)$ is a dually dense sublattice of D_2 .

Note that in Boolean algebras, dense and dually dense conditions are equivalent; in distributive lattices this equivalence does not hold (see [9] for some known characterizations of density and dual density in distributive lattices).

For the case of contact algebras [24] and distributive contact lattices [9] we introduced the notion of *C-separability* as follows. Let $D_1 \leq D_2$; we say that D_1 is a *C-separable sublattice of* D_2 if the following condition is satisfied:

(C-separable) $(\forall a_2, b_2 \in D_2)(a_2\overline{C}b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_2 \leq a_1, b_2 \leq b_1, a_1\overline{C}b_1))$. For the case of EDC-lattices we modified this notion adding two additional clauses corresponding to the relations \widehat{C} and \ll just having in mind the definitions of these relations in contact algebras. Namely

Definition 5.4 C-separability. Let $D_1 \preceq D_2$; we say that D_1 is a C-separable EDC-sublattice of D_2 if the following conditions are satisfied:

 $(C ext{-}separability for }C)$ -

$$(\forall a_2, b_2 \in D_2)(a_2\overline{C}b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_2 \le a_1, b_2 \le b_1, a_1\overline{C}b_1)).$$

 $(C ext{-separability for }\widehat{C})$ -

$$(\forall a_2, b_2 \in D_2)(a_2\overline{C}b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_2 + a_1 = 1, b_2 + b_1 = 1, a_1\overline{C}b_1)).$$

 $(C\text{-}separability for \ll)$

$$(\forall a_2, b_2 \in D_2)(a_2 \ll b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_2 \leq a_1, b_2 + b_1 = 1, a_1 \overline{C}b_1)).$$

If h is an embedding of the lattice D_1 into the lattice D_2 then we say that h is a C-separable embedding if the sublattice $h(D_1)$ is a C-separable sublattice of D_2 .

The notion of a C-separable embedding h is defined similarly. The following lemma is analogous to a similar result from [24] (Theorem 2.2.2) and from [9] (Lemma 5).

Lemma 5.5 Let D_1 , D_2 be EDC-lattices and D_1 be a C-separable EDC-sublattice of D_2 . Then:

- (i) If D_1 is a dually dense EDC-sublattice of D_2 , then D_1 satisfies the axiom (Ext C) iff D_2 satisfies the axiom (Ext C),
 - (ii) D_1 satisfies the axiom (Con C) iff D_2 satisfies the axiom (Con C),
 - (iii) D_1 satisfies the axiom (Nor 1) iff D_2 satisfies the axiom (Nor 1),
- (iv) D_1 satisfies the axiom (U-rich \ll) iff D_2 satisfies the axiom (U-rich \ll),
 - (v) D_1 satisfies the axiom (U-rich \widehat{C}) iff D_2 satisfies the axiom (U-rich \widehat{C}).

Proof. Conditions (i), (ii) and (iii) have the same proof as in Theorem 2.2.2 from [24].

- (iv) (\Rightarrow) Suppose that D_1 satisfies the axiom (U-rich \ll), $a_2, b_2 \in D_2$ and let $a_2 \ll b_2$. Then by (C-separability for \ll) we obtain: $(\exists a_1, b_1 \in D_1)(a_2 \leq a_1, b_2 + b_1 = 1, a_1\overline{C}b_1)$. Since D_1 is a sublattice of D_2 then $a_1, b_1 \in D_2$. From $a_2 \leq a_1$ and $a_1\overline{C}b_1$ we get $a_2\overline{C}b_1$. Thus we have just proved: $(a_2 \ll b_2 \to (\exists b_1 \in D_2)(b_2 + b_1 = 1 \text{ and } a_2\overline{C}b_1)$ which shows that D_2 satisfies (U-rich \ll).
- (\Leftarrow) Suppose that D_2 satisfies the axiom (U-rich \ll), $a_1,b_1 \in D_1$ (hence $a_1,b_1 \in D_2$) and let $a_1 \ll b_1$. Then by (U-rich \ll) for D_2 we get: $(\exists c_2 \in D_2)(b_1 + c_2 = 1, a_1\overline{C}c_2)$. Since $a_1, c_2 \in D_2$ and $a_1\overline{C}c_2$, then by (C-separability for C) we get: $(\exists a'_1, b'_1 \in D_1)(a_1 \leq a'_1, c_2 \leq b'_1, a'_1\overline{C}b'_1)$. Combining the above results we get: $1 = b_1 + c_2 \leq b_1 + b'_1$ and $a_1\overline{C}b'_1$. We have just proved the following: $a_1 \ll b_1 \to (\exists b'_1 \in D_1)(b_1 + b'_1 = 1, a_1\overline{C}b'_1)$ which shows that D_1 satisfies (U-rich \ll).
 - (v) The proof is similar to that of (iv). \square .

The notion of \widehat{C} -separable sublattice can be defined in a dual way as follows:

Definition 5.6 Suppose that $D_1 \leq D_2$; we say that D_1 is a \widehat{C} -separable EDC-sublattice of D_2 if the following condition is satisfied:

 $(\widehat{C}$ -separability for C) -

$$(\forall a_2, b_2 \in D_2)(a_2Cb_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_1 + a_2 = 1, b_1 + b_2 = 1, a_1\overline{\widehat{C}}b_1)),$$

 $(\widehat{C}\operatorname{-separability} for \widehat{C})$ -

$$(\forall a_2, b_2 \in D_2)(a_2\widehat{C}b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_1 \le a_2, b_1 \le b_2, a_1\overline{\widehat{C}}b_1)),$$

 $(\widehat{C}\operatorname{-separability} for \ll)$ -

$$(\forall a_2, b_2 \in D_2)(a_2 \ll b_2 \Rightarrow (\exists a_1, b_1 \in D_1)(a_1 + a_2 = 1, b_1 \le b_2, a_1 \overline{\widehat{C}}b_1)).$$

The notion of a \widehat{C} -separable embedding h is defined as in definition 5.4.

The following lemma is dual to Lemma 5.5 and can be proved in a dual way.

Lemma 5.7 Let D_1, D_2 be EDC-lattices and D_1 be a \widehat{C} -separable EDC-sublattice of D_2 ; then:

- (i) If D_1 is a dense EDC-sublattice of D_2 , then D_1 satisfies the axiom (Ext \widehat{C}) iff D_2 satisfies the axiom (Ext \widehat{C}),
 - (ii) D_1 satisfies the axiom (Con \widehat{C}) iff D_2 satisfies the axiom (Con \widehat{C}),
 - (iii) D_1 satisfies the axiom (Nor 2) iff D_2 satisfies the axiom (Nor 2).
- (iv) D_1 satisfies the axiom (O-rich \ll) iff D_2 satisfies the axiom (O-rich \ll).
 - (v) D_1 satisfies the axiom (O-rich \widehat{C}) iff D_2 satisfies the axiom (O-rich \widehat{C}).

Corollary 5.8 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice and $\underline{B} = (B, C)$ be a contact algebra. Then:

- (i) If h is a C-separable embedding of \underline{D} into \underline{B} then \underline{D} must satisfy the axioms (U-rich $\ll)$ and (U-rich $\widehat{C})$.
- (ii) If h is a \hat{C} -separable embedding of \underline{D} into \underline{B} then \underline{D} must satisfy the axioms (O-rich \ll) and (O-rich \hat{C}).
- **Proof.** (i) Note that by Lemma 5.2 \underline{B} satisfies the axioms (U-rich \ll) and (U-rich \widehat{C}). Then by Lemma 5.5 (iv) and (v) \underline{D} satisfies the axioms (U-rich \ll) and (U-rich \widehat{C}).
 - (ii) Similarly to (i) the proof follows from Lemma 5.2 and Lemma 5.7. \square

PART II: TOPOLOGICAL REPRESENTATIONS OF EXTENDED DISTRIBUTIVE CONTACT LATTICES

The aim of this second part of the paper is to investigate several kinds of topological representations of EDC-lattices. We concentrate our attention mainly on topological representations with some "good properties" in the sense of Section 5.1: dual density and C-separability, and their dual versions - density and \hat{C} -separability.

6 Topological models of EDC-lattices

We assume some familiarity of the reader with the basic theory of topological spaces: (see [12]). First we recall some notions from topology. By a topological space we mean a set X provided with a family C(X) of subsets, called closed sets, which contains the empty set \varnothing , the whole set X, and is closed with respect to finite unions and arbitrary intersections. Fixing C(X) we say that X is endowed with a topology. A subset $a \subseteq X$ is called open if it is the complement of a closed set. A family of closed sets $\mathbf{CB}(X)$ is called a closed basis of the topology if every closed set can be represented as an intersection of sets from $\mathbf{CB}(X)$. In a similar way the topology of X can be characterized by the family O(X) of open sets: it contains the empty set, X and is closed under finite intersections and arbitrary unions. A family $\mathbf{OB}(X)$ of open sets is called an open basis of the topology if every open set can be represented as an union of sets from $\mathbf{OB}(X)$. X is called semiregular space if it has a closed base of regular closed sets or an open base of regular open sets.

We remaind the definitions of two important topological operations on sets - closure operation Cl, and interior operation Int. Namely Cl(a) is the intersection of all closed sets of X containing a and Int(a) is the union of all open sets included in a. Note that the operations Cl and Int are interdefinable: Cl(a) = -Int(-a) and Int(a) = -Cl(-a). Using the bases CB(X) and OB(X) the definitions of closure and interior operations have the following useful expressions:

```
x \in Cl(a) iff (\forall b \in \mathbf{CB}(X))(a \subseteq b \to x \in b), x \in Int(a) iff (\exists b \in \mathbf{OB}(X))(b \subseteq a \text{ and } x \in b).
```

We say that a is a regular closed set if a = Cl(Int(a)) and a is a regular open set if a = Int(Cl(a)). It is a well known fact that the set RC(X) of all regular closed subsets of X is a Boolean algebra with respect to the relations, operations and constants defined as follows: $a \le b$ iff $a \subseteq b$, $0 = \emptyset$, 1 = X, $a+b=a \cup b$, $a \cdot b = Cl(Int(a \cap b)$, $a^* = Cl(-a)$ where $-a = X \setminus a$. If we define a contact C by aCb iff $a \cap b \ne \emptyset$ then we obtain the standard topological model of contact algebra.

Another topological model of contact algebra is by the set RO(X) of regular open subsets of X. The relevant definitions are as follows: $a \leq b$ iff $a \subseteq b$, $0 = \emptyset$, 1 = X, $a \cdot b = a \cap b$, $a + b = Int(Cl(a \cup b), a^* = Int - a$. The contact relation is aCb iff $Cl(a) \cap Cl(b) \neq \emptyset$.

Note that these two models are isomorphic.

Topological model of EDC-lattice by regular-closed sets. Consider the contact algebra RC(X) of regular closed subsets of X. Let us remove the operation a^* and define the relations \widehat{C} and \ll topologically according to their definitions in contact algebra as follows:

```
a\widehat{C}b iff Cl(-a) \cap Cl(-b) \neq \emptyset iff (equivalently) Int(a) \cup Int(b) \neq X. a \ll b iff a \cap Cl(-b) = \emptyset iff (equivalently) a \subseteq Int(b).
```

Obviously the obtained structure is a model of EDC-lattice. Also any distributive sublattice of RC(X) with the same definitions of the relations C, \widehat{C} and \ll is a model of EDC-lattice. These models are considered as standard topological models of EDC-lattice by regular closed sets.

Topological model of EDC-lattice by regular-open sets. Consider the contact algebra RO(X) of regular open subsets of X. Let us remove the operation a^* from the contact algebra RO(X) and define the relations \widehat{C} and \ll topologically according to their definitions in the contact algebra as follows:

 $a\widehat{C}b$ iff $Cl(Int(-a) \cap Cl(Int(-b)) \neq \emptyset$ iff (equivalently) $a \cup b \neq X$,

 $a \ll b$ iff $Cl(a) \cap Cl(Int(-b)) = \emptyset$ iff (equivalently) $Cl(a) \subseteq b$.

Obviously the obtained structure is another standard topological model of EDC-lattice and any distributive sublattice of RO(X) with the same relations C, \hat{C} and \ll is also a model of EDC-lattice.

The main aim of PART II of the paper is the topological representation theory of EDC-lattices related to the above two standard models. The first simple result is the following representation theorem.

Theorem 6.1 Topological representation theorem for EDC-lattices. Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Then:

- (i) There exists a topological space X and an embedding of \underline{D} into the contact algebra RC(X) of regular closed subsets of X.
- (ii) There exists a topological space Y and an embedding of \underline{D} into the contact algebra RO(Y) of regular open subsets of Y.

Proof. It is shown in [6] that every contact algebra is isomorphic to a subalgebra of the contact algebra RC(X) of regular closed subsets of some topological space X, and dually, that it is also isomorphic to a subalgebra of the contact algebra RO(Y) of the regular open subsets of some topological space Y. Then the proof follows directly from this result and the Corollary $3.12.\square$

The above theorem is not the best one, because it can not be extended straightforwardly to EDC-lattices satisfying some of the additional axioms mentioned in Section 5. That is why we will study in the next sections representation theorems based on embeddings satisfying some of the good conditions described in Section 5.1. Before going on let us remaind some other topological facts, which will be used later on.

A topological space X is called:

- normal if every pair of closed disjoint sets can be separated by a pair of open sets;
- κ -normal [19] if every pair of regular closed disjoint sets can be separated by a pair of open sets;
- weakly regular [10] if it is semiregular and for each nonempty open set a there exits a nonempty open set b such that $Cl(a) \subseteq b$;
- *connected* if it can not be represented by a sum of two disjoint nonempty open sets;
- T_0 if for every pair of distinct points there is an open set containing one of them and not containing the other; X is called T_1 if every one-point set is a closed set, and X is called Hausdorff (or T_2) if each pair of distinct points can be separated by a pair of disjoint open sets.
- compact if it satisfies the following condition: let $\{A_i : i \in I\}$ be a non-empty family of closed sets of X such that for every finite subset $J \subseteq I$ the intersection $\bigcap \{A_i : i \in J\} \neq \emptyset$, then $\bigcap \{A_i : i \in I\} \neq \emptyset$.

The following lemma relates topological properties to the properties of the relations C, \hat{C} and \ll and shows the importance of the additional axioms for EDC-lattices.

Lemma 6.2 (i) If X is semiregular, then X is weakly regular iff RC(X) satisfies any of the axioms (Ext C), (Ext \widehat{C}).

- (ii) X is κ -normal iff RC(X) satisfies any of the axioms (Nor 1), (Nor 2) and (Nor 3).
- (iii) X is connected iff RC(X) satisfies any of the axioms (Con C), (Con \widehat{C}).
- (iv) If X is compact and Hausdorff, then RC(X) satisfies (Ext C), (Ext \widehat{C}) and (Nor 1), (Nor 2) and (Nor 3).
- **Proof.** A variant of the above lemma concerning only axioms (Ext C), (Nor 1) and (Con C) was proved, for instance, in [10]. Having in mind the equivalence of some of the mentioned axioms in RC(X), it is obvious that the present formulation is equivalent to the cited result from [10].

6.1 Looking for good topological representations of EDC-lattices

The following topological theorem proved in [9] (Theorem 4) gives necessary and sufficient conditions for a closed base of a topology to be semiregular.

Theorem 6.3 First characterization theorem for semiregularity.

Let X be a topological space and let $\mathbf{CB}(X)$ be a closed basis for X. Suppose that "·" is a binary operation defined on the set $\mathbf{CB}(X)$ such that $(\mathbf{CB}(X), \varnothing, X, \cup, \cdot)$ is a lattice. Then:

- 1. The following conditions are equivalent:
 - (a) CB(X) is U-extensional.
 - (b) $\mathbf{CB}(X) \subseteq RC(X)$.
 - (c) For all $a, b \in \mathbf{CB}(X)$, $a \cdot b = Cl(Int(a \cap b))$.
 - (d) $(\mathbf{CB}(X), \varnothing, X, \cup, \cdot)$ is a dually dense sublattice of the Boolean algebra RC(X).
- 2. If any of the (equivalent) conditions (a),(b),(c) or (d) of 1. is fulfilled then:
 - (a) $(CB(X), \emptyset, X, \cup, \cdot)$ is a *U*-extensional distributive lattice.
 - (b) X is a semiregular space.

The following is a corollary of the above theorem.

Corollary 6.4 [9] Let X be a topological space, let $L = (L, 0, 1, +, \cdot)$ be a lattice and let h be an embedding of the upper semi-lattice (L, 0, 1, +) into the lattice C(X) of closed sets of X. Suppose that the set $\mathbf{CB}(X) = \{h(a) : a \in L\}$ forms a closed basis for the topology of X. Then:

- 1. The following conditions are equivalent:
 - (a) L is U-extensional.
 - (b) $\mathbf{CB}(X) \subseteq RC(X)$.
 - (c) For all $a, b \in L$, $h(a \cdot b) = Cl(Int(h(a) \cap h(b)))$.

- (d) h is a dually dense embedding of L into the Boolean algebra RC(X).
- 2. If any of the (equivalent) conditions (a),(b),(c) or (d) of 1. is fulfilled then:
 - (a) L is a U-extensional distributive lattice.
 - (b) X is a semiregular space.

A dual version of Theorem 6.3 is the following one.

Theorem 6.5 Second characterization theorem for semiregularity. Let X be a topological space and let $\mathbf{OB}(X)$ be an open basis for X. Suppose that + is a binary operation defined on the set $\mathbf{OB}(X)$ such that $(\mathbf{OB}(X), \varnothing, X, \cap, +)$ is a lattice. Then:

- 1. The following conditions are equivalent:
 - (a) OB(X) is O-extensional.
 - (b) $\mathbf{OB}(X) \subseteq RO(X)$.
 - (c) For all $a, b \in \mathbf{OB}(X)$, $a + b = Int(Cl(a \cup b))$.
 - (d) $(\mathbf{OB}(X), \varnothing, X, \cap, +)$ is a dually dense sublattice of the Boolean algebra RO(X).
- 2. If any of the (equivalent) conditions (a),(b),(c) or (d) of 1. is fulfilled then:
 - (a) $(\mathbf{OB}(X), \varnothing, X, \cap, +)$ is an O-extensional distributive lattice.
 - (b) X is a semiregular space.

The following is a corollary of the above theorem.

Corollary 6.6 Let X be a topological space, let $L = (L,0,1,+,\cdot)$ be a lattice and let h be an embedding of the lower semi-lattice $(L,0,1,\cdot)$ into the lattice O(X) of open sets of X. Suppose that the set $\mathbf{OB}(X) = \{h(a) : a \in L\}$ forms an open basis for the topology of X. Then:

- 1. The following conditions are equivalent:
 - (a) L is O-extensional.
 - (b) $\mathbf{OB}(X) \subseteq RO(X)$.
 - (c) For all $a, b \in L$, $h(a + b) = Int(Cl(h(a) \cup h(b)))$.
 - (d) h is a dense embedding of L into the Boolean algebra RO(X).
- 2. If any of the (equivalent) conditions (a),(b),(c) or (d) of 1. is fulfilled then:
 - (a) L is a O-extensional distributive lattice.
 - (b) X is a semiregular space.

Remark 6.7 (i) Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Corollary 6.4 shows that if we want to represent \underline{D} by a dually dense embedding h into the contact algebra RC(X) of some topological space X such that the topology of X to be determined by the set $\mathbf{CB}(X) = \{h(a) : a \in D\}$ considered as a closed base for X we must require that the lattice \underline{D} is U-extensional, i.e. to satisfy the axiom (Ext \widehat{O}) (extensionality of underlap). If in addition we want to apply the

good properties of Lemma 5.5 then we must assume that h is also a C-separable embedding into RC(X). But then Corollary 5.8 implies that \underline{D} must satisfy also the axioms (U-rich \ll) and (U-rich \widehat{C}).

(ii) Similar to the above conclusion is the following. Corollary 6.6 shows that if we want to represent \underline{D} by a dense embedding h into the contact algebra RO(X) of some topological space X such that the topology of X to be determined by the set $\mathbf{OB}(X) = \{h(a) : a \in D\}$ considered as an open base for X we must require that the lattice \underline{D} is O-extensional, i.e. to satisfy the axiom (Ext O) (extensionality of overlap). If in addition we want to apply the good properties of Lemma 5.7 then we must assume that h is also a \widehat{C} -separable embedding into RO(X). But then Corollary 5.8 implies that \underline{D} must satisfy also the axioms (O-rich \ll) and (O-rich \widehat{C}). \square

Definition 6.8 U-rich and O-rich EDC-lattices. Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Then:

- (i) \underline{D} is called U-rich EDC-lattice if it satisfies the axioms (Ext \widehat{O}), (U-rich \ll) and (U-rich \widehat{C}).
- (ii) \underline{D} is called O-rich EDC-lattice if it satisfies the axioms (Ext O), (O-rich \ll) and (O-rich \widehat{C}).

The aim of the next sections is to develop the topological representation theory of U-rich and O-rich EDC-lattices.

7 Topological representation theory of U-rich EDC-lattices

The aim of this section is to develop a topological representation theory for U-rich EDC-latices. According to Theorem 6.3 we will look for a representation with regular closed sets. To realize this we will follow the representation theory of contact algebras by regular closed sets developed in [6, 24], updating the results of Section 4 from [9] to the case of U-rich EDC-lattices. We will consider also extensions of U-rich EDC-lattices with some of the additional axioms mentioned in Section 5. The scheme of the representation procedure is the following: for each U-rich EDC-lattice D from a given class, determined by the additional axioms, we will do the following:

- Define a set X(D) of "abstract points" of D,
- define a topology in X(D) by the set $\mathbf{CB}(X(D)) = \{h(a) : a \in D\}$, considered as a closed base of the topology, where h is the intended embedding of Stone type: $h(a) = \{\Gamma : \Gamma \text{ is "abstract point" and } a \in \Gamma\}$. X(D) is called the canonical topological space of D and h is called canonical embedding,
- establish that h is a dual dense embedding of the lattice \underline{D} into the Boolean algebra RC(X(D)) of regular closed sets of the space X(D).

We will consider separately the cases of representations in T_0 , T_1 and T_2 spaces which requires introducing different "abstract points".

7.1 Representations in T_0 spaces

Troughout this section we consider that $\underline{D} = (D, C, \widehat{C}, \ll)$ is a U-rich EDC-lattice.

Abstract points of \underline{D} .

As in [9], we consider the abstract points of \underline{D} to be clans (see [6] for the origin of this notion). The definition is the following. A subset $\Gamma \subseteq D$ is a *clan* if it satisfies the following conditions:

(Clan 1) $1 \in \Gamma$, $0 \notin \Gamma$,

(Clan 2) If $a \in \Gamma$ and $a \leq b$, then $b \in \Gamma$,

(Clan 3) If $a + b \in \Gamma$, then $a \in \Gamma$ or $b \in \Gamma$,

(Clan 4) If $a, b \in \Gamma$ then aCb.

 Γ is a maximal clan if it is maximal with respect to the set-inclusion. We denote by CLAN(D) (MaxCLAN(D)) the set of all (maximal) clans of \underline{D} .

The notion of clan is an abstraction from the following natural example. Let X be a topological space and RC(X) be the contact algebra of regular-closed subsets of X and let $x \in X$. Then the set $\Gamma_x = \{a \in RC(X) : x \in a\}$ is a clan.

Now we will present a construction of clans which is similar to the constructions of clans in contact algebras. First we will introduce a new canonical relation between prime filters.

Definition 7.1 Let U, V be prime filters. Define a new canonical relation R_C (R_C -canonical relation) between prime filters as follows:

$$UR_CV \leftrightarrow_{def} (\forall a \in U)(\forall b \in V)(aCb).$$

Let us note that the relation R_C depends only on C and can be defined also for filters. It is different from the canonical relation between prime filters defined in Section 3.3, but the presence of U-rich axioms makes it equivalent to R^c as it can be seen from the following lemma.

Lemma 7.2 (i) R_C is reflexive and symmetric relation.

(ii) If \underline{D} satisfies the axioms (U-rich \ll) and (U-rich \widehat{C}) then $R_C = R^c$.

Proof. (i) follows from the axioms (C4) and (C5).

(ii) The inclusion $R^c \subseteq R_C$ follows directly by the definition of R^c . For the converse inclusion suppose UR_CV . To show UR^cV we have to inspect the four cases of the definition of R^c .

Claim 1: $a \in U$ and $b \in V$ implies aCb. This is just by the definition of R_C .

Claim 2: $a \in U$ and $b \notin V$ implies $a \not\ll b$. For the sake of contradiction suppose $a \in U$ and $b \notin V$ but $a \ll b$. Then by axiom (U-rich \ll) ($a \ll b \to (\exists c)(b+c=1)$ and $a\overline{C}c$), we obtain b+c=1 and $a\overline{C}c$. Conditions b+c=1 and $b \notin V$ imply $c \in V$. But $a \in U$, so aCc - a contradiction.

Claim 3: $a \notin U$ and $b \in V$ implies $b \not\ll a$. The proof is similar to the proof of Claim 2.

Claim 4: $a \notin U$ and $b \notin V$ implies $a\widehat{C}b$. The proof is similar to the proof of Claim 2 by the use of axiom

(U-rich \widehat{C}) $a\overline{\widehat{C}}b \to (\exists c,d)(a+c=1,b+c=1 \text{ and } c\overline{C}d)$. \square

The following statement lists some facts about the relation R_C .

Facts 7.3 [7, 6, 9].

- 1. Let F,G be filters and FR_CG then there are prime filters U,V such that $F\subseteq U$, $G\subseteq V$ and UR_CV .
- 2. For all $a, b \in D$: aCb iff there exist prime filters U, V such that UR_CV , $a \in U$ and $b \in V$.

In the following lemma we list some facts about clans (see, for instance, [6, 9]).

Facts 7.4 1. Every prime filter is a clan.

- 2. The complement of every clan is an ideal.
- 3. If Γ is a clan and F is a filter such that $F \subseteq \Gamma$, then there is a prime filter U such that $F \subseteq U \subseteq \Gamma$. In particular, if $a \in \Gamma$, then there exists a prime filter U such that $a \in U \subseteq \Gamma$.
- 4. Every clan Γ is the union of all prime filters contained in Γ .
- 5. Every clan is contained in a maximal clan.
- 6. Let Σ be a nonempty set of prime filters such that for every $U, V \in \Sigma$ we have UR_CV and let Γ be the union of the elements of Σ . Then Γ is a clan and every clan can be obtained in this way.
- 7. Let U, V be prime filters, Γ be a clan and $U, V \subseteq \Gamma$. Then UR_CV and UR^cV .

Lemma 7.5 Let Γ be a clan and $a \in D$. Then the following two conditions are equivalent:

- (i) $(\forall c \in D)(a+c=1 \rightarrow c \in \Gamma)$,
- (ii) There exists a prime filter $U \subseteq \Gamma$ such that $a \notin U$.

Proof. (i) \rightarrow (ii). Suppose that (i) holds. It is easy to see that the set $F = \{c : a+c=1\}$ is a filter. The complement $\overline{\Gamma}$ of Γ is an ideal (Facts 7.4) and hence $\overline{\Gamma} \oplus (a]$ is an ideal. We will show that $F \cap \overline{\Gamma} \oplus (a] = \emptyset$. Suppose the contrary. Then there is a c such that a+c=1 (and hence by (i) $c \in \Gamma$) and $c \in \overline{\Gamma} \oplus (a]$. Then there is $x \in \overline{\Gamma}$ such that $c \leq x+a$. From here we get: $1 = a+c \leq a+x+a = x+a$, hence x+a=1 and by (i) $-x \in \Gamma$, contrary to $x \in \overline{\Gamma}$. Now we can apply Filter-extension Lemma and obtain a prime filter U extending F such that $U \cap \overline{\Gamma} \oplus (a] = \emptyset$. It follows from here that $a \notin U$, $U \cap \overline{\Gamma} = \emptyset$ which implies $U \subseteq \Gamma$.

(ii) \rightarrow (i). Suppose (ii) holds: $U \subseteq \Gamma$ and $a \notin U$. Suppose a+c=1. Then $c \in U \subseteq \Gamma$, so $c \in \Gamma$ - (i) is fulfilled. \square

Defining the canonical topological space X(D) of \underline{D} and the canonical embedding h.

Define the Stone like embedding: $h(a) = \{\Gamma \in CLAN(D) : a \in \Gamma\}$ and consider the set $\mathbf{CB}(X) = \{h(a) : a \in D\}$ as a closed base of the topology in X(D) = CLAN(D).

Lemma 7.6 The space X(D) is semiregular and h is a dually dense embedding of \underline{D} into the contact Boolean algebra RC(X(D)).

Proof. Using the properties of clans, one can easily check that $h(0) = \emptyset$, h(1) = X, and that $h(a+b) = h(a) \cup h(b)$. This shows that the set $\mathbf{CB}(X(D)) = \{h(a) : a \in D\}$ is closed under finite unions and, in fact, it is a closed basis for the topology of X. Also we have the implication: $a \leq b$ then $h(a) \subseteq h(b)$.

To show that h is an embedding we use the fact that prime filters are clans and prove that $a \not\leq b$ implies $h(a) \not\subseteq h(b)$. Indeed, from $a \not\leq b$ it follows by the theory of distributive lattices (see [2]) that there exists a prime filter U (which is also a clan) such that $a \in U$ (so $U \in h(a)$) and $b \notin U$ (so, $U \notin h(b)$), which proves that $h(a) \not\subseteq h(b)$. Consequently, h is an embedding of the upper semi-lattice (D,0,1,+) into the lattice of closed sets of the space X(D). By Corollary 6.4, X(D) is a semiregular space and h is a dually dense embedding of D into the Boolean algebra RC(X). It remains to show that h preserves the relations C, \widehat{C} and \ll . This follows from the following claim.

Claim 7.7 (i) Let Γ be a clan and $a \in D$. Then following equivalence holds: $\Gamma \in h(a)$ iff there exists a prime filter U such that $a \in U \subseteq \Gamma$.

(ii) Let Γ be a clan and $a \in D$. Then following conditions are equivalent:

```
(I) (\forall c \in D)(a+c=1 \rightarrow c \in \Gamma),

(II) \Gamma \in Cl(-h(a)),
```

(III) There exists a prime filter U such that $a \notin U \subseteq \Gamma$.

```
(iii) aCb iff h(a) \cap h(b) \neq \emptyset,

(iv) a \not\ll b iff h(a) \cap Cl(-h(b)) \neq \emptyset.

(v) a\widehat{C}b iff Cl(-h(a)) \cap Cl(-h(b)) \neq \emptyset,
```

Proof of the claim. (i) follows easily from Facts 7.4 (3.).

(ii) The proof of $(I) \leftrightarrow (II)$ follows by the following sequence of equivalences:

```
 (\forall c \in D)(a+c=1 \to c \in \Gamma) \text{ iff} 
 (\forall c \in D)(h(a) \cup h(c) = X(D) \to \Gamma \in h(c)) \text{ iff} 
 (\forall c \in D)(-h(a) \subseteq h(c) \to \Gamma \in h(c)) \text{ iff} 
 \Gamma \in Cl(-h(a))
```

The first equivalence holds because h is an embedding of the upper semilattice (D,0,1,+) into the lattice of closed sets of the space X(D), the third equivalence uses the fact that the set $\{h(c): c \in D\}$ is a closed base of the topology of X(D).

The equivalence $(I) \leftrightarrow (III)$ is just the Lemma 7.5.

- (iii) (\Rightarrow) Suppose aCb, then by Lemma 3.8 (i) there exist prime filters U, and V such that UR^cV , $a \in U$ and $b \in V$. Let $\Gamma = U \cup V$. By Facts 7.4 Γ is a clan, obviously containing a and b, which implies $h(a) \cap h(b) \neq \emptyset$.
- (\Leftarrow) Suppose $h(a) \cap h(b) \neq \emptyset$. Then there exists a clan Γ containing a and b, hence aCb.
- (iv) (\Rightarrow) Suppose $a \not\ll b$. Then by Lemma 3.8 (ii) there exist prime filters U, V such that UR^cV , $a \in U$ and $b \notin V$. Let $\Gamma = U \cup V$, then Γ is a clan

containing U and V. So, $a \in \Gamma$ and hence $\Gamma \in h(a)$. From the condition $b \notin V \subseteq \Gamma$ we obtain by (ii) that $\Gamma \in Cl(-h(b))$ and hence $h(a) \cap Cl(-h(b)) \neq \emptyset$.

- (\Leftarrow) Suppose $h(a) \cap Cl(-h(b)) \neq \varnothing$. Then there exists a clan $\Gamma \in h(a)$ and $\Gamma \in Cl(-h(b))$. It follows by (i) that there exists a prime filter U such that $a \in U \subseteq \Gamma$ and by (ii) we obtain that there exists a prime filter V such that $b \notin V \subseteq \Gamma$. Condition $U, V \subseteq \Gamma$ implies by Facts 7.4 (7.) that UR^cV . Using the properties of the relation R^c and $a \in U$ and $b \notin V$ we get $a \not\ll b$.
- (v) The proof of (v) is similar to the proof of (iv) with the use of Lemma 3.9. This finishes the proof of Lemma 7.6 \square

Lemma 7.8 The following conditions are true for the canonical space X(D):

- (i) X(D) is T_0 .
- (ii) X(D) is compact.

Proof. The proof is the same as the proof of Lemma 19 from [9].

Lemma 7.9 The mapping h is a C-separable embedding of D into RC(X(D)).

Proof. This lemma was proved in [9] by a special construction. Since the definition of C-separability for EDC-lattices uses an extended definition for which the special construction from [9] does not hold, in this paper we give a new proof deducing the statement from the compactness of the space X(D).

We have to proof the following three statements, corresponding to the three clauses of the condition of C-separability (see Definition 5.4).

(C-separability for C) $(\forall \alpha, \beta \in RC(X(D)))(\alpha \cap \beta = \emptyset \rightarrow (\exists a, b \in D)(\alpha \subseteq h(a), \beta \subseteq h(b), h(a) \cap h(b)) = \emptyset$.

(C-separability for \widehat{C}) $(\forall \alpha, \beta \in RC(X(D))(Cl(-\alpha) \cap Cl(-\beta) = \emptyset \rightarrow (\exists a, b \in D)(\alpha \cup h(a) = X(D), \beta \cup h(b) = X(D), h(a) \cap h(b) = \emptyset)$.

(C-separability for \ll) $(\forall \alpha, \beta \in RC(X(D))(\alpha \cap Cl(-\beta) = \emptyset \rightarrow (\exists a, b \in D)(\alpha \subseteq h(a), \beta \cup h(b) = X(D), h(a) \cap h(b) = \emptyset).$

As an example we shall prove the condition (C-separability for C). The proofs for the other two conditions are similar.

Proof of (C-separability for C). Let $\alpha, \beta \in RC(X(D))$ and $\alpha \cap \beta = \emptyset$. Since α and β are closed sets they can be represented as intersections from the elements of the basis $\mathbf{CB}(X(D)) = \{h(c) : c \in D\}$ of X(D). So there are subsets $A, B \subseteq \mathbf{CB}(X(D))$ such that $\alpha = \bigcap \{h(c) : h(c) \in A\}$ and $\beta = \bigcap \{h(c) : h(c) \in B\}$. Then $\alpha \cap \beta = \bigcap \{h(c) : h(c) \in A\} \cap \bigcap \{h(c) : h(c) \in B\} = \emptyset$. By the compactness of X(D) (Lemma 7.8 (ii)), there are finite subsets $A_0 \subseteq A$ and $B_0 \subseteq B$ such that $\alpha \cap \beta = \bigcap \{h(c) : h(c) \in A_0\} \cap \bigcap \{h(c) : h(c) \in B_0\} = \emptyset$. Let $A_0 = \{h(c_1), ..., h(c_n)\}$ and $B_0 = \{h(d_1), ..., h(d_m)\}$ and let $a = c_1 \cdot ... \cdot c_n$ and $b = d_1 \cdot ... \cdot d_m$. Then $h(a) \subseteq h(c_i)$, i = 1...n and from here we get $h(a) \subseteq h(c_1) \cap ... \cap h(c_n)$. Analogously we obtain that $h(b) \subseteq h(d_1) \cap ... \cap h(d_m)$. Consequently $h(a) \cap h(b) \subseteq (h(c_1) \cap ... \cap h(c_n) \cap (h(d_1) \cap ... \cap h(d_m)) = \emptyset$, so $h(a) \cap h(b) = \emptyset$. Also we have $\alpha \subseteq h(c)$ for all $h(c) \in A$ and consequently for all $h(c) \in A_0$. Hence $\alpha \subseteq h(c_1) \cdot ... \cdot h(c_n) = h(c_1 \cdot ... \cdot c_n) = h(a)$, so $\alpha \subseteq h(a)$. Analogously we get $\beta \subseteq h(b)$. \square

The following theorem is the main result of this section.

Theorem 7.10 Topological representation theorem for U-rich EDC-lattices

Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an U-rich EDC-lattice. Then there exists a compact semiregular T_0 -space X and a dually dense and C-separable embedding h of \underline{D} into the Boolean contact algebra RC(X) of the regular closed sets of X. Moreover:

- (i) \underline{D} satisfies (Ext C) iff RC(X) satisfies (Ext C); in this case X is weakly regular.
- (ii) \underline{D} satisfies (Con C) iff RC(X) satisfies (Con C); in this case X is connected.
- (iii) \underline{D} satisfies (Nor 1) iff RC(X) satisfies (Nor 1); in this case X is κ -normal.

Proof. Let X be the canonical space X(D) of \underline{D} and h be the canonical embedding of \underline{D} . Then, the theorem is a corollary of Lemma 7.6, Lemma 7.8, Lemma 7.9 and Lemma 6.2.

Note that Theorem 7.10 generalizes several results from [6, 10] to the distributive case.

7.2 Representations in T_1 spaces

The aim of this section is to obtain representations of some U-rich EDC-lattices in T_1 -spaces extending the corresponding results from [9]. The constructions will be slight modifications of the corresponding constructions from the previous section, so we will be sketchy.

Let $\underline{D}=(D,C,\widehat{C},\ll)$ be an U-rich EDC-lattice. In the previous section the abstract points were clans and this guarantees that the representation space is T_0 . To obtain representations in T_1 spaces we assume abstract points to be maximal clans, so for the canonical space of \underline{D} we put X(D)=MaxCLAN(D) and define the canonical embedding h to be $h(a)=\{\Gamma\in MaxCLAN(D): a\in\Gamma\}$. The topology in X(D) is defined considering the set $\mathbf{CB}(X(D))=\{h(a): a\in D\}$ to be the closed base for the space. Note that in general, without additional axioms we can not prove in this case that h is an embedding. In order to guarantee this we will assume that \underline{D} satisfies additionally the axiom of C-extensionality

(Ext C)
$$a \neq 1 \rightarrow (\exists b \neq 0)(a\overline{C}b)$$
.

Note that in this case, due to U-extensionality (see Section 5), the lattice \underline{D} satisfies also the axiom

(EXT C)
$$a \leq b \rightarrow (\exists c)(aCc \text{ and } b\overline{C}c)$$
,

which is essential in the proof that h is an embedding.

Lemma 7.11 The space X(D) is a semiregular and h is a dually dense embedding of D into the contact Boolean algebra RC(X(D)).

Proof. The proof is similar to the proof of Lemma 7.6, so we will indicate only the differences. First we show that h is an embedding of the upper semi-lattice (D,0,1,+) into the lattice of closed sets of the space X(D). The only new thing which we have to show is: If $a \not\leq b$ then $h(a) \not\subseteq h(b)$. To do this suppose $a \not\leq b$. Then by axiom (EXT C) there exists $c \in D$ such that aCc but $b\overline{C}c$. Condition aCc implies that there exist prime filters U,V such that UR^cV , $a \in U$ and $c \in V$. Let $\Gamma_0 = U \cup V$. Γ_0 is a clan and by Facts 7.4 it is contained in a maximal clan Γ . Obviously $a, c \in \Gamma$, so $\Gamma \in h(a)$. But $b\overline{C}c$ implies that

 $b \notin \Gamma$ (otherwise we will get bCc). Conditions $\Gamma \in h(a)$ and $\Gamma \notin h(b)$ show that $h(a) \not\subseteq h(b)$. Thus, by Corollary 6.4 h is a dually dense embedding of \underline{D} into the Boolean algebra RC(X(D)). It remains to show that h preserves the relations C, \widehat{C} and \ll . The proof is almost the same as in the corresponding proof of Lemma 7.6. The only new thing is when we construct a certain clan from prime filters satisfying the relation UR^cV in the form $U \cup V$, then we extend it into a maximal clan. Note also that Claim 7.7 remains true. We demonstrate this by considering only the preservation of \ll . We have to show:

```
a \not\ll b \text{ iff } h(a) \cap Cl(-h(b) \neq \emptyset
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- (\Rightarrow) Suppose $a \not\ll b$. Then by Lemma 3.8 $(\exists U, V \in PF(D))(a \in U \text{ and } b \not\in V \text{ and } UR^cV)$. Define $\Gamma_0 = U \cup V$. Γ_0 is a clan containing U and V. Extend Γ_0 into a maximal clan Γ . Then Γ contains a, so $\Gamma \in h(a)$. We have also that $b \not\in V \subseteq \Gamma$, so by the Claim 7.7 $\Gamma \in Cl(-h(b))$.
 - (\Leftarrow) The proof is identical to the corresponding proof from Lemma 7.6. \square

Lemma 7.12 The space X(D) satisfies the following conditions:

- (i) X(D) is T_1 ,
- (ii) X(D) is compact,
- (iii) h is C-separable embedding.
- **Proof.** (i) Let Γ be an arbitrary maximal clan. The space X(D) is T_1 iff the singleton set $\{\Gamma\}$ is closed, i.e. $Cl(\{\Gamma\}) = \{\Gamma\}$. This follows by the maximality of Γ as follows. Let Δ be a maximal clan. Then:
- $\Delta \in Cl(\{\Gamma\}) \text{ iff } (\forall c \in D)(\{\Gamma\} \subseteq h(c) \to \Delta \in h(c)) \text{ iff } (\forall c \in D)(\Gamma \in h(c) \to \Delta \in h(c)) \text{ iff } (\forall c \in D)(c \in \Gamma \to c \in \Delta \text{ iff } \Gamma \subseteq \Delta) \text{ iff } \Gamma = \Delta \text{ iff } \Delta \in \{\Gamma\}.$

This chain shows that indeed $Cl(\{\Gamma\}) = \{\Gamma\}.$

- (ii) The proof is similar to the proof of Lemma 7.8 (ii)
- (iii) follows from (ii) as in the proof of Lemma 7.9. \square

Theorem 7.13 Topological representation theorem for C-extensional U-rich EDC-lattices Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be a C-extensional U-rich EDC-lattice. Then there exists a compact weakly regular T_1 -space X and a dually dense and C-separable embedding h of \underline{D} into the Boolean contact algebra RC(X) of the regular closed sets of X. Moreover:

- (i) \underline{D} satisfies (Con C) iff RC(X) satisfies (Con C); in this case X is connected.
- (ii) \underline{D} satisfies (Nor 1) iff RC(X) satisfies (Nor 1); in this case X is κ -normal.

Proof. The proof follows from Lemma 7.11, Lemma 7.12 and Lemma 6.2.

7.3 Representations in T_2 spaces

In the previous section we proved representability in T_1 spaces of U-rich EDC-lattices satisfying the axiom of C-extensionality (Ext C). The T_1 property of the topological space was guaranteed by the fact that abstract points are maximal clans. In this section we will show that adding the axiom (Nor 1) we can obtain representability in compact T_2 -spaces. The reason for this is that the axiom (Nor 1) makes possible to use new abstract points - the so called clusters, which are maximal clans satisfying some additional properties yielding T_2 separability of

the topological space. Clusters have been used in the compactification theory of proximity spaces (see more about their origin in [23]). They have been adapted in algebraic form in the representation theory of contact algebras in [6, 26]. In [9] their definition and some constructions are modified for the distributive case. We remaind below the corresponding definition.

Definition 7.14 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. A clan Γ in \underline{D} is called a cluster if it satisfies the following condition:

(Cluster) If for all $b \in \Gamma$ we have aCb, then $a \in \Gamma$. We denote the set of clusters in \underline{D} by CLUSTER(D).

Let us note that not in all EDC-lattices there are clusters. The following lemma shows that the axiom (Nor 1) guarantees existence of clusters and some important properties needed for the representation theorem.

Lemma 7.15 [9] Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. Then:

- (i) Every cluster is a maximal clan.
- (ii) If \underline{D} satisfies (Nor 1) then every maximal clan is a cluster.
- (iii) If Γ and Δ are clusters such that $\Gamma \neq \Delta$, then there are $a \notin \Gamma$ and $b \notin \Delta$ such that a + b = 1.

To build the canonical space X(D) we assume in this section that $\underline{D} = (D, C, \widehat{C}, \ll)$ is an U-rich EDC-lattice satisfying the axioms (Ext C) and (Nor 1). We define X(D) = CLUSTER(D), $h(a) = \{\Gamma \in CLUSTER(D) : a \in \Gamma\}$ and define the topology in X(D) considering the set $\mathbf{CB}(X) = \{h(a) : a \in D\}$ as a basis for closed sets in X(D). Since the points of X(D) are maximal clans, just as in Section 7.2 we can prove the following lemma.

Lemma 7.16 The space X(D) is a semiregular and h is a dually dense embedding of D into the contact Boolean algebra RC(X(D)).

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Lemma 7.17 (i) X(D) is T_2, (ii) X(D) is compact,
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- (iii) h is C-separable embedding.
- **Proof.** (i) To show that the space X(D) is T_2 suppose that Γ, Δ are two different clusters. We have to find two disjoint open sets A, B such that $\Gamma \in A$ and $\Delta \in B$. By Lemma 7.15 (iii) there are $a, b \in D$ such that $a \notin \Gamma$ and $b \notin \Delta$ such that a + b = 1. Then by Lemma 7.16 we get $\Gamma \notin h(a), \Delta \notin h(b)$ and $h(a) \cup h(b) = X(D)$, hence $-h(a) \cap -h(b) = \emptyset$. Define A = -h(a), B = -h(b). Since h(a) and h(b) are closed sets, then A and B are open sets which separate the abstract points Γ and Δ .

The proof of (ii) and (iii) is the same as the proof of (ii) and (iii) in Lemma 7.12. \Box

Theorem 7.18 Topological representation theorem for U-rich EDC-lattices satisfying (Ext C) and (Nor 1). Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an U-rich EDC-lattice satisfying (Ext C) and (Nor 1). Then there exists a compact T_2 -space X and a dually dense and C-separable embedding h of \underline{D} into the Boolean contact algebra RC(X) of the regular closed sets of X. Moreover \underline{D} satisfies (Con C) iff RC(X) satisfies (Con C) and in this case X is connected.

Proof. The proof follows from Lemma 7.16, Lemma 7.17 and 6.2. \square Let us note that this theorem generalizes several theorems from [6, 24, 27, 26]

8 Topological representation theory of O-rich EDC-lattices

This section is devoted to the theory of dense representations for O-rich EDC-latices (see Definition 6.8). According to Theorem 6.5 we will look for dense representations with regular open sets. This case is completely dual to the corresponding theory developed in Section 7. For this reason we will only sketch the main representation scheme and the definitions of abstract points for the T_0 , T_1 and T_2 representations.

The representation scheme is dual to the scheme presented in Section 7:

- Define a set X(D) of "abstract points" of \underline{D} ,
- define a topology in X(D) by the set $\mathbf{OB}(X(D)) = \{h(a) : a \in D\}$, considered as an open base of the topology, where h is the intended embedding of Stone type: $h(a) = \{\Gamma : \Gamma \text{ is "abstract point" and } a \in \Gamma\}$. X(D) is called the canonical topological space of \underline{D} and h is called canonical embedding,
- establish that h is a dense embedding of the lattice \underline{D} into the Boolean algebra RO(X(D)) of regular open sets of the space X(D).

For the case of T_0 dense representation we consider a notion of abstract point which is dual to the notion of clan. This is the so called E-filter (Efremovich filter). E-filters were used in the theory of proximity spaces (see [23]). In the context of contact algebras they were introduced for the first time in [6]. The definition adapted for the language of EDC-lattices is the following.

Definition 8.1 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an EDC-lattice. A subset $\Gamma \subseteq D$ is called an E-filter if it satisfies the following properties:

```
(E-fil 1) \Gamma is a proper filter in \underline{D}, i.e. 0 \notin \Gamma, (E-fil 2) If a \notin \Gamma and b \notin \Gamma, then a\widehat{C}b.
```

 Γ is a minimal E-filter if it is minimal in the set of all E-filters of \underline{D} with respect to set inclusion.

This definition comes as an abstraction from the following natural example. Let X be a topological space, $x \in X$ and RO(X) be the set of all regular-open sets of X. Then the set $\Gamma_x = \{a \in RO(X) : x \in a\}$ is an E-filter in the contact algebra RO(X). Note that the definition of E-filter is based not on the relation of contact C, but on the dual contact \widehat{C} .

A general construction of E-filters can be obtain dualizing the construction of clans from Section 7.1. Just to show how this dual construction goes on and how the O-rich axioms works, we will repeat some steps omitting the proofs.

First we will introduce a new canonical relation between prime filters.

Definition 8.2 Let U, V be prime ideals. Define a new canonical relation $\widehat{R}_{\widehat{C}}$ ($\widehat{R}_{\widehat{C}}$ -canonical relation) between prime ideals as follows:

$$U\widehat{R}_{\widehat{C}}V \leftrightarrow_{def} (\forall a \in U)(\forall b \in V)(a\widehat{C}b).$$

If U, V are prime filters then we define $UR_{\widehat{C}}V \leftrightarrow_{def} \overline{U}\widehat{R}_{\widehat{C}}\overline{V}$.

Let us note that the relation $\widehat{R}_{\widehat{C}}$ depend only on \widehat{C} and can be defined also for ideals. It is different from the canonical relation \widehat{R}^c between prime ideals defined in Section 3.3, but the presence of O-rich axioms makes it equivalent to \widehat{R}^c as it is stated in the following lemma.

Lemma 8.3 (i) $\widehat{R}_{\widehat{C}}$ is a reflexive and symmetric relation.

(ii) If \underline{D} satisfies the axioms (O-rich \ll) and (O-rich \widehat{C}), then $\widehat{R}_{\widehat{C}} = \widehat{R}^c$.

The following statement lists some facts about the relation R_C .

- **Facts 8.4** 1. Let F, G be ideals and $F\widehat{R}_{\widehat{C}}G$ then there are prime ideals U, V such that $F \subseteq U$, $G \subseteq V$ and $U\widehat{R}_{\widehat{C}}V$.
 - 2. For all $a, b \in D$: $a\widehat{C}b$ iff there exist prime ideals U, V such that $U\widehat{R}_{\widehat{C}}V$, $a \in U$ and $b \in V$.
 - 3. For all $a,b \in D$: $a\widehat{C}b$ iff there exist prime filters U,V such that $UR_{\widehat{C}}V$, $a \notin U$ and $b \notin V$.

In the following lemma we list some facts about E-filters.

Facts 8.5 1. Every prime filter is an E-filter.

- 2. If Γ is an E-filter and $a \notin \Gamma$, then there exists a prime filter U such that $\Gamma \subseteq U$ and $a \notin U$.
- 3. Every E-filter Γ is the intersection of all prime filters containing Γ .
- 4. Every E-filter contains a minimal E-filter.
- 5. Let Σ be a nonempty set of prime filters such that for every $U, V \in \Sigma$ we have $UR_{\widehat{C}}V$ and let Γ be the intersection of the elements of Σ . Then Γ is an E-filter and every E-filter can be obtained in this way.
- 6. Let U,V be prime filters, Γ be an E-filter, $\Gamma\subseteq U$ and $\Gamma\subseteq V$. Then $UR_{\widehat{C}}V$ and $UR^{c}V$.

Using the above facts one can prove the following representation theorem.

Theorem 8.6 Let $\underline{D} = (D, C, \widehat{C}, \ll)$ be an O-rich EDC-lattice. Then there exists a compact semi-regular space X and a dense and \widehat{C} -separable embedding h from \underline{D} into the contact algebra RO(X) of regular-open sets of X. Moreover:

- (i) If \underline{D} satisfies (Ext \widehat{C}), then X is weakly regular,
- (ii) If \underline{D} satisfies (Con \widehat{C}), then X is a connected space,
- (iii) If \underline{D} satisfies (Nor 2), then X is κ -normal.

Abstract points for dense representations in T_1 spaces are minimal E-filters and abstract points for dense representations in T_2 spaces are duals of clusters introduced in [6] under the name *co-clusters*. We adapt this notion for the language of EDC-lattices as follows:

Definition 8.7 An E-filter Γ is called co-cluster if it satisfies the following condition:

(Co-cluster) If $(\forall b \notin \Gamma)(a\widehat{C}b)$, then $a \notin \Gamma$. (or, equivalently, if $a \in \Gamma$, then $(\exists b \notin \Gamma)(a\overline{\widehat{C}b})$).

Let us show, for instance, the following statement for co-clusters, which is dual to the corresponding property for clusters as maximal clans:

Lemma 8.8 Every co-cluster is a minimal E-filter.

Proof. Suppose that Γ is a co-cluster which is not a minimal E-filter. Then there exists an E-filter Δ such that $\Delta \subset \Gamma$, so $a \in \Gamma$ and $a \notin \Delta$ for some a. Then there exists $b \notin \Gamma$ such that $a\overline{\widehat{C}}b$. From here we get $b \in \Delta$. Consequently $b \in \Gamma$ - a contradiction. \square

We left to the reader to formulate and proof the dual analog of Theorem 7.13 and Theorem 7.18.

9 Concluding remarks

In this paper we generalized the notion of contact algebra by weakening the algebraic part to distributive lattice. One solution of this problem was given in [9] including in the definition only the contact relation. However, the obtained axiomatization in [9] is in a sense "incomplete", because it does not contain the definable in the Boolean case mereotopological relations of dual contact \widehat{C} and non-tangential inclusion \ll and its dual \gg and in this sense the system is not closed under duality. We succeed in this paper to axiomatize all these relations considered as primitives on the base of distributive lattices by means of universal first-order axioms. The resulting system is called "extended distributive contact lattice" (EDC-lattice). In this way we obtain, among others, the following two results. First, EDC-lattice is closed under duality, and second, it can be considered as an axiomatization of the universal fragment of contact algebras in the language of distributive lattices with the relations C, \widehat{C} and \ll . We developed topological representation theory of EDC-lattices by means of regular closed and regular open sets generalizing in a quite non-trivial way the corresponding representation theory for contact algebras. Considering this representation theory on a weaker lattice base provided a deeper insight into the interaction of some notions taking place in the representation, which cannot be seen in the Boolean case. For instance we show the role of extensionality of underlap and overlap relations in case of dual dense and dense embeddings.

Our future plans include building of new logics for qualitative spatial representation and reasoning based on EDC-lattices, studying the standard logical problems related to them: axiomatizability, decidability or undecidability, complexity. A good source for possible generalizations and extensions is the paper [3] containing many examples of spatial logics based on contact and precontact algebras.

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